

Formal Methods in Software Engineering

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Logic Problem

- In order to better understand the logical rules and the logical connectors, let us look at a logic problem.
- On an island there are two types of residents: Knights and knaves.
- Knights always speak truth and knaves always tell a lie.
- Suppose a stranger arrives there and asks someone “Are you a knight”.
- The person replies, “If I am a knight then I will eat my hat.”

Logic Problem

- Now we have to proof that the person will eat his hat.
- Using the logical rules, we have to proof that our hypothesis is true.
- First step is to convert the problem into propositions and see how we can model/formulate this problem.
- Lets use the symbol A to denote the proposition "A is a knight."
- Let H be the proposition, "A eats his hat".

Logic Problem

- A is a knight: A
- A eats his hat: H
- If I am a knight then I'll eat my hat:
 - $A \Rightarrow H$

Logic Problem

- The person X is equivalent to the statement S, “If I am a knight then I will eat my hat.”
- $X \Leftrightarrow S$
- Person X is equivalent to the statement because, if X is knight the statement is true, if X is knave the statement is false.
- So X and S both have the same truth value.
- So if person A is making the statement $A \Rightarrow H$, we will have
- $A \Leftrightarrow A \Rightarrow H$

Logic Problem

- We have seen that $(X \Leftrightarrow S)$
- Therefore $(A \Leftrightarrow A \Rightarrow H)$
- Now our objective is to logically deduce H.

Logic Problem

- First we will try to proof it using truth table.

A	H	$A \Rightarrow H$	$A \Leftrightarrow (A \Rightarrow H)$
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Truth Table Columns

Proof Using Truth Table

A	H	$A \Rightarrow H$
T	T	T
F	T	T
T	F	F
F	F	T

Proof Using Truth Table

A	$A \Rightarrow H$	$A \Leftrightarrow (A \Rightarrow H)$
T	T	T
F	T	F
T	F	F
F	T	F

Proof Using Truth Table

A	H	$A \Rightarrow H$	$A \Leftrightarrow (A \Rightarrow H)$
T	T	T	T
F	T	T	F
T	F	F	F
F	F	T	F

- Hence there is only one valid row and in that valid row, the value of H is true, means the proposition H, "A eats his hat" is true.
- Hence we conclude that A has to eat his eat.

Problems with Truth Table

- The problem with truth tables is that you can make a truth table if number of propositions are small.
- But if number of variables are large, it is not easy to make a truth table.
- For example, if we have 2 variables, we will have $2^2 = 4$ rows in the truth table.
- But what if we have 10 variables, then we will have $2^{10} = 1K$ rows.
- It won't be possible to make a truth table.
- So for large variables, we can use deduction and equivalence.

Proof using equivalences

- $A \Leftrightarrow (A \Rightarrow H)$
 - $\equiv A \Leftrightarrow (\text{not } A \text{ or } H)$
 - $\equiv (A \text{ and } (\text{not } A \text{ or } H)) \text{ or}$
 - $(\text{not } A \text{ and } \text{not } (\text{not } A \text{ or } H))$

Proof using equivalences

- $A \text{ and } (\text{not } A \text{ or } H)$ //Apply Distributive law
- $\equiv (A \text{ and not } A) \text{ or } (A \text{ and } H)$
- $\equiv \text{false or } (A \text{ and } H)$
- $\equiv A \text{ and } H$

Proof using equivalences

- $\text{not } A \text{ and not } (\text{not } A \text{ or } H)$
 $\equiv \text{not } A \text{ and } (A \text{ and not } H)$
 $\equiv (\text{not } A \text{ and } A) \text{ and not } H$
 $\equiv \text{false and not } H$
 $\equiv \text{false}$

//Apply De-Morgans Law

//Apply Associative Law

Proof using equivalences

- Hence

$$A \Leftrightarrow (\text{not } A \text{ or } H)$$

$$\equiv (A \text{ and } H) \text{ or false}$$

$$\equiv A \text{ and } H$$

- Therefore $A \Leftrightarrow A \Rightarrow H$ is equivalent to $A \text{ and } H$.
- Hence our outcome is $A \text{ and } H$ i.e. A is a knight and H (A has to eat his hat).

Rules of Inference and Logical Deduction

- There are two category of rules that we will use for logical deduction and proofs.
 - Introduction
 - Elimination

Rules of Inference and Logical Deduction

- This is the first rule of Introduction.
- Given two propositions p , q , that are true.
- We can conclude that p and q is true.

Introduction

$$\frac{p \quad q}{p \text{ and } q}$$

Rules of Inference and Logical Deduction

- This is the second rule of Introduction.
- If we are given only one proposition, say only p or only q.
- If p is given, means p is true. We can say that p or q is true.
- Similarly if q is given i.e. q is true. We can say that p or q is true.

Introduction	
$\frac{p}{p \text{ or } q}$	$\frac{q}{p \text{ or } q}$

Rules of Inference and Logical Deduction

- Another important rule of introduction is, If we have q and q is true. We can say $p \Rightarrow q$.
- As we have seen in implication truth table, if q is true, $p \Rightarrow q$ will always be true.
- No matter what is the value of p .

Introduction
$\frac{[p] \quad q}{p \Rightarrow q}$

Rules of Inference and Logical Deduction

- So first rule of Elimination is: given p and q , means p and q is true. We can conclude p is true.
- Similarly, given p and q , means p and q is true. We can conclude q is true.
- So we can eliminate either p or q .
- Because if P and q is true, means p as well as q is true independent of each other.

Elimination	
$\frac{p \text{ and } q}{p}$	$\frac{p \text{ and } q}{q}$

Rules of Inference and Logical Deduction

- This is the second rule of Elimination.
- If we are given that p is true and also $p \Rightarrow q$ is true.
- We can conclude that q must be true.

Elimination	
p	$p \Rightarrow q$
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q	

Rules of Inference and Logical Deduction

- Another important rule of elimination is, if we are given p and not p .
- What does it mean? It means we have contradiction.
- We have two contradictory statements. So the outcome will be false.
- Similarly if we start from False i.e. if we have a false argument or assumption.
- From a false assumption , we can conclude anything.

Elimination		
p	not p	false
<hr/>		<hr/>
false		p

End of Slides