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Q 1:-  $(64)_{10} = (?)_2$

$$\begin{array}{r|l} 2 & 64 \\ \hline 2 & 32-0 \\ \hline 2 & 16-0 \\ \hline 2 & 8-0 \\ \hline 2 & 4-0 \\ \hline 2 & 2-0 \\ \hline 1 & 1-0 \end{array}$$

$$(64)_{10} = (1000000)_2$$

Q 1:- (b)  $(0111111)_2 = (?)_{10}$

$$(0 \times 2^7) + (1 \times 2^6) + (1 \times 2^5) + (1 \times 2^4) + (1 \times 2^3) + (1 \times 2^2) + (1 \times 2^1) + (1 \times 2^0)$$

$$= 0 + 64 + 32 + 16 + 8 + 4 + 2 + 1$$

$$= (127)_{10} \text{ Ans.}$$

Q 1:- (c) ~~4DE~~  $(4D7F)_{16} = (?)_{10}$

$$= 4 \times 16^3 + D \times 16^2 + 7 \times 16^1 + F \times 16^0$$

$$= 4 \times 16^3 + 13 \times 16^2 + 7 \times 16^1 + 15 \times 16^0$$

$$= 4 \times 4096 + 13 \times 256 + 7 \times 16 + 15$$

$$= 16384 + 3328 + 112 + 15$$

$$= 19839$$

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Q4:- (d)  $(128)_{10} = (?)_{16}$

$(80)_{16}$

$$\begin{array}{r|l} 16 & 28 \\ \hline & 8 - 0 \end{array}$$

Q4:- (e)  $(3A6F) = (?)_2$

3      A      6      F  
0011    1010    0110    1111

$(0011101001101111)_2$

Q4:- (f)  $(110000111100101)_2 = (?)_{16}$

↓      ↓      ↓      ↓  
12    3      14    5

$(C3E5)_{16}$  Ans.

Q4:- (g)  $(11111111)_2 = \pm (?)_{10}$

MSB is 1 shows number is negative so:

$$\begin{aligned} & -1 \times 2^7 + (1 \times 2^6) + (1 \times 2^5) + (1 \times 2^4) + (1 \times 2^3) \\ & + (1 \times 2^2) + (1 \times 2^1) + (1 \times 2^0) \end{aligned}$$

$$= -1 \times 128 + (1 \times 64) + (1 \times 32) + (1 \times 16)$$

$$+(1 \times 8) + (1 \times 4) + (1 \times 2) + (1 \times 1)$$

$$-128 + 127$$

$$= -1 \text{ Ans.}$$

Q1:- (h)  $(-16)_{10} = (?)_2$

$$(00010000)_2$$

Taking 2's complement

$$\begin{array}{r|l} 2 & 16 \\ \hline 2 & 8 - 0 \\ \hline 2 & 4 - 0 \\ \hline 2 & 2 - 0 \\ \hline & 1 - 0 \end{array}$$

$$00010000$$

$$11101111$$

+ 1

2's complement

$$\underline{11110000}$$

$$(11110000)_2 \text{ Ans.}$$

Q1:- (i)  $(01111111)_2 - (00000111)_2$

~~1's complement~~

~~$$\begin{array}{r} 01111111 \\ 11111000 \end{array}$$~~

$$01111111$$

$$00000111$$

$$\underline{01111000}$$

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2's complement of (01111000)

$$\begin{array}{r} \phantom{000} \\ \phantom{000} \\ \phantom{000} \\ 10000111 \\ + 1 \\ \hline \underline{10001000} \end{array}$$

(10001000)<sub>2</sub> Ans.

Q 1: - (j)  $6D_{16} - 3F_{16}$

$$\begin{array}{r} \swarrow \quad \downarrow \quad \downarrow \quad \searrow \\ \underline{0110} \quad \underline{1101} \quad \underline{0011} \quad \underline{1111} \end{array}$$

$$6D = 0110 \ 1101$$

$$3F = 0011 \ 1111$$

Taking 1<sup>st</sup> number of second number

$$01111111$$

$$10000000$$

Now add both numbers

$$1101101$$

$$\underline{1000000}$$

$$\leftarrow \textcircled{1} 0101101$$

Carry

add carry in result

$$0101101$$

+ 1

$$\underline{0101110}$$

(2E)<sub>16</sub> Ans.

Q2:- Write short note on each of the following:

(a) Embedded systems:-

Embedded system is a combination of computer sys software and hardware which is either fixed in capability or programmable. An embedded system can be either an independent system, or it can it can be a part of a large system. It is mostly designed for a specific function or functions within a larger system.

(b) Device driver:-

Most commonly known as a driver, a device driver or a hardware driver is a group of files that enable one or more hardware devices to communicate with the computer's operating system. Without drivers, the computer would not be able to

to hardware devices, such as a printer.

(c) Virtual machine concept:-

A virtual machine (VM) is a software program or operating system that not only exhibits the behaviour of a separate computer, but is also capable of performing tasks such as running applications and programs like a separate computer. A virtual machine, usually known as a guest is created within another computing environment referred as a "host".

(d) Instruction execution cycle:-

Instruction execution means a program to be executed by a processor consists of a set of instructions stored in memory.

In instruction execution cycle the time period during which one instruction is fetched from memory and execute when computer given an instruction in machine language.

- Each instruction is further divided into sequence of phases.
- After the execution of program counter is incremented, to point to the next instructions.

(e) Motherboard chipset :-

A chipset is a group of interdependent motherboard chips or integrated circuits that control the flow of data and instructions between the central processing unit (CPU) or microprocessor and external devices.

A chipset controls external buses, memory and some peripherals. A CPU unable to function without impeccable chipset timing.

(f) Access levels for input-output operations:-

1) High level operations

→ Is more flexible and usually more convenient.

→ Hides complexity from the programming.

→ When using the I/O functions high-level I/O is slower as compared to the low-level I/O.

2) Low-level operations:

→ provides direct access to files and devices.

→ Is complex (Buffer management is to be done by the programmer)

→ When using I/O functions, low-level I/O is faster as compared to the high-level I/O.

→ Uses a "file descriptor" to track the status of the file.



(9) Basic parts of an assembly language instruction

An assembly instruction has 4 basic parts:

- 1) Label
- 2) Mnemonic
- 3) Operands
- 4) Comments

### Labels:-

There are two types of labels:

- Data labels
- Code labels

### Data labels:-

Data labels are just like variables in any other language. For example, C language instruction: `int myVar10;` Similarly, in assembly language we use data labels as variables which hold values for us. For example `count` is a variable in the following instruction

```
count DWORD 100
```

Data label declare in the data segment code.

### Code Label:-

Code labels are just the names which are used by other instructions to jump from particular position to that position where data label is located. It must end with a semicolon(;).

For example:

target:

mov ax, bx

...

jmp target

target is a data label, which is called by the jmp instruction.

### Mnemonic:-

A small word that acts as an identifier for the instructions. The mnemonics are written in code segment. In following examples mov, sub, add, jmp, call, and mul are the mnemonic:

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- mov  $\Rightarrow$  Move/assign one value to another
- sub  $\Rightarrow$  Subtract one value from another
- add  $\Rightarrow$  adds two values
- jmp  $\Rightarrow$  Jump to a specific location
- Call  $\Rightarrow$  Call a procedure/module
- mul  $\Rightarrow$  Multiply two values.

### Operands:-

Operand is a remaining part of the instruction. There can be 0 to 3 operands. Where each operand is either a memory operand, a register, an input/output port, or a constant expression.

Following are the operands types:

- Constant expression
- Constant
- memory
- register

### Comments:-

You should have used comments in your programs before. As you know there are two types of comments :- 1. Singl-line and

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2. Multi-line comments. In assembly single-line comments begin with ; (semicolon) whereas, multi-line comments begin with a comment directive and symbolic/character selected by the programmer itself and also end with the same symbol/character.

Q3: (a) Differentiate Between each of following:

(a) Assembly language and High-level language :-

1) → In assembly language programs written for one processor will not run on another type of processor.

→ In high level language programs runs independently of processor type

2) Performance and accuracy of assembly language code are better than a high-level.

→ High-level languages have to give extra instructions to run code on computer.

(3) Code of assembly languages ~~have~~ is difficult to understand and debug than a high-level

→ One or two statements of high-level languages into many assembly languages codes.

4) Assembly language can communicate better than a high-level language. Some type of hardware actions can only be performed by assembly language.

→ In assembly language we can directly read pointers at a physical address which is not possible in high-level.

(b) Protected mode and real address mode:-

The 80286 introduced something into the x86 architecture called "protected mode".

protected mode differed from the original mode of the 8086, which was later dubbed "real mode", in that areas of memory could be physically isolated by the processor itself to prevent illegal writes to other programs could be running in memory at the same time, but any program could access

any area of memory and, therefore, if malicious or errant, for example, could take down the entire system.

### (C) Assembler and linker

The main output produced by assembler an input assembly language source file is the translation of the file into an object file. The object files produced by the assembler are relocatable files that hold code and/or data. They are input files for linker.

### (d) Code label and data label :-

Data label is the label that we use to define data, as we defined memory locations num1, num2... etc in our program.

Code label is the label that we have on code as we see in case of conditional jump (Label1) and is normally used for loop control statement.

(f) Line comment and block comment

→ A single line comment and as applied as only applied sigle to a single line in the "source code"

→ A block comment usually refers to a paragraph of text. A block comment has a start symbol and an end symbol and everything between is ignored by the computer.

(g) Equal-sign directive and EQU directive :-

Equal-sign

→ Expression is 32-bit integer

→ May be redefined

→ name is called a symbolic constant

Syntax :-

name = expression

EQU

Define a symbol as either an integer or text.

Cannot be redefined

Syntax :-

name EQU expression ; integer expression

name EQU symbol ; existing symbol name

name EQU <text> ; any text



Q4: (a) Explain the concept of portability as it applies to programming languages.

Ans: A language whose source program can be compiled and run on a wide variety of computer system is said to be portable.

Q4: (b) Why would a high-level language not be an ideal tool for writing a program that directly accesses a particular brand of printer?

Ans: A high-level language may not be provide for direct hardware access. Even it does awkward coding technique must often be used, resulting in possible maintain problem.

Q4: (c) Why was unicode invented?

Ans: Unicode is a universal computing standard to represent texts in most writing system. It was invented to store most of the

words character. it is stated during 1987.

Q4:- (d) If  $W = 11101100$ ,  $X = 00010011$ , and  $Y = 00111100$ , then find  $Z = W \vee X \wedge \neg Y$ .

Ans:

$$Z = W \vee X \wedge \neg Y$$

$$Y = 00111100$$

$$\neg Y = 11000011$$

$$\Rightarrow X \wedge \neg Y = \begin{array}{r} 00010011 \\ 11000011 \\ \hline \end{array}$$

$$\text{AND } \underline{\underline{00000011}}$$

$$W \vee X \wedge \neg Y$$

$$11101100$$

$$00000011$$

OR

$$\underline{\underline{11101111}}$$

Q4:- (e) Create a truth table to show all possible inputs and outputs for the Boolean function described by  $\neg(A \vee B)$

$$\neg(A \vee B)$$

A	B	$A \vee B$	$\neg(A \vee B)$
F	F	F	T
F	T	T	F
T	F	T	F
T	T	T	F

Q4:- (f) Why does memory access take more machine cycles than register access?

Ans: Conventional memory is outside the CPU, and it responds more slowly to access requests. Registers are hard-wired inside the CPU.

Q4:- (g) Discuss the basic program execution registers used in x86 32-bit processors.

Ans:-

XMM Registers:-

The x86 architecture also contain eight 128-bit register called streaming SIMD extension to the instruction.

MMX Registers:-

MMX technology improves the performance of intel processor when implementing advanced multimedia and communication application. The 64-bit MMX register support special instruction called SIMD (single-instruction) Multiple data.

Segment Register:-

In real address mode, 16-bit segment register indicate base addresses of preassigned memory areas named segments.

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Basic programming Execution

Register :-

Registers are high-speed storage location inside the CPU designed to be accessed at much higher speed than conventional memory when a processing loop is optimized for speed.

Q 28:- Discuss the following `stack` directives.

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Stack:-

The `stack` directive tells how many bytes of memory to reserve for the runtime stack.

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4096 happen to correspond to the size of a memory page in this processor system for managing memory.

### MODEL:-

This tells the assembler which memory model to use in 32-bit program. We use the flat memory model which is associated with the processor's protected model.

### 386:-

```
-> model flat  
-> stack 4096  
Exit process PROTO  
dw Exit code: DWORD
```

The 386 directive identifies it as a 32-bit program. Line 2 uses the flat memory model and window requires the model conversion to be used.

Line 3 set aside 4096 bytes of storage.

Lin 4 declares a proto type for the Exit process function.

CODE:-

code  
main PROC

it is the beginning of the code area of the program (meaning what's after-word is usually the main procedure.

Data:-

The DATA directive creates a near data segment.

→ This segment contain the frequency used data for your program

→ Data segment can occupy

\* up to 64k in MS-DOS

\* or upto 512 mega-byte under flat model in window NT

PROTO:-

The proto directives by using the invoke directives

Syntax:-

Label PROTO [[distance]][[language-type]]  
[[parameter]]:tag...]]



parameter :-

distance (32-bit MASM only)  
(optional) used in 16-bit memory model to override the default and indicate NEAR or FAR calls

PROC :-

MASK starts and ends of a procedure block called label.

The statements in the block can be called with the call instruction or INVOKE directives

ENDP :-

Mark the end of procedure name previously begun with PROC.

END :-

Directives for END of file command. That's end of file here as you are using "main" you have to end it with.

Q6: (a) Write a program that calculates the following expression:  $A = (A+B) - (C+D)$

Ans:  $A = (A+B) - (C+D)$

data ; data segment, read and write

Var A BYTE 10

Var B BYTE 20

Var C BYTE 30

Var D BYTE 40

final var BYTE

code ; code segment, read-only

main PROC

MOV al, Var A

MOV bl, Var B

MOV cl, Var C

MOV dl, Var D

ADD al, bl ; compute (A+B)

ADD cl, dl ; compute (C+D)

SUB al, cl ; compute (A+B) - (C+D)

MOV final\_val, al

call Dump Regs

exit  
main EBP  
EBP main

Q6:- (b) Show the order of individual bytes in memory for the following doubleword variable using little endian order:

Ans: When placed in memory at offset 0000, 78h would be stored in the first byte 56h would be stored in the second byte and the remain bytes would be at offset 0002 and 0003 as show in figure :

1 2 3 4 5 6 7 8h

0000	78
0001	56
0002	34
0003	12

Q6:- (C)

data

string byte "Assembly language is easy", 0  
string size byte ?

code

```
mov eax size EOF string 1  
mov string size, eax
```

Q6:- (D) Let the Arithmetic operation

$$is \ x = -a + (b - c)$$

```
INCLUDE vine 32.inc
```

• data

```
x DWORD ? ; uninitialized variable  
a DWORD 10 ; initialize variable 'a'  
b DWORD 20 ; initialize variable 'b'  
c DWORD 15 ; initialize variable 'c'
```

• Code

```
main PROC
```

```
mov eax, a ; move value of 'a' into 'eax'
```

```
neg eax ; EAX = -10
```

```
mov ebx, b ; move value of b into 'ebx'
```

```
sub ebx, c ; EBX = 11
```

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```
add eax, ebx ; perform -10 + 11  
mov x, eax ; 1
```

```
call DUMPRegs
```

```
exit
```

```
main ENDP
```

```
END main
```