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I.D : 15385

Semester : 3rd

Exam : Mid Assignment

Subject : DLD

Q1, Convert each of the following:

a, $45.25_{10} = (?)_2$

For the whole decimal numbers, repeated division by 2 will be used.

2	45	
2	22	1
2	11	0
2	5	1
2	2	1
	1	0

Now for the fractional part, we will use repeated multiplication by 2.

$$0.25 \times 2 = 0.50 \text{ --- } 0$$

$$0.50 \times 2 = 1.00 \text{ --- } 1$$

hence,

$$(45.25)_{10} = (101101.01)_2$$

Answer

$$b/ \quad 10000000.1010_2 = (?)_{10}$$

We will use weighted notation to convert the given binary number to decimal.

$$\begin{aligned} &= 1 \times 2^7 + 1 \times 2^1 + 0 \times 2^{-2} + 1 \times 2^{-3} + 0 \times 2^{-4} \\ &= 1 \times 128 + 0.5 + 0.125 \\ &= (128.625)_{10} \end{aligned}$$

$$(10000000.1010)_2 = (128.625)_{10}$$

Answer

$$c/ \quad 4D7F_{16} = (?)_{10}$$

Using weighted notation.

$$\begin{aligned} &= 4 \times 16^3 + 13 \times 16^2 + 7 \times 16^1 + 15 \times 16^0 \\ &= 16384 + 3328 + 112 + 15 \\ &= (19839)_{10} \end{aligned}$$

Answer

$$\text{dr } 128_{10} = (?)_{16}$$

Using repeated division by 16.

$$\begin{array}{r|l} 16 & 128 \\ \hline 16 & 8 \quad 0 \end{array}$$

$$(128)_{10} = (80)_{16}$$

Answer

$$\text{e, } 3A6F_{16} = (?)_2$$

Using Hexa-Binary table

<u>3</u>	<u>A</u>	<u>6</u>	<u>F</u>
0011	1010	0110	1111

So,

$$\{ \cancel{(3A6F)}_{16} = \cancel{(11101001101111)} \} \times$$

$$\textcircled{(3A6F)}_{16} = (11101001101111)_2$$

Answer

$$f. \quad 110000111100101_2 = (?)_{16}$$

making groups of four

$$\begin{array}{cccc} \underline{1100} & \underline{0011} & \underline{1110} & \underline{0101} \\ C & 3 & E & 5 \end{array}$$

so,

$$(110000111100101)_2 = (C3E5)_{16}$$

Ans

$$g. \quad 6173_8 = (?)_{10}$$

Using weighted notation

$$= 6 \times 8^3 + 1 \times 8^2 + 7 \times 8^1 + 3 \times 8^0$$

$$= 3072 + 64 + 56 + 3$$

$$= (3195)_{10}$$

so,

$$(6173)_8 = (3195)_{10}$$

Answer.

h/ $169_{10} = (?)_8$

Using repeated division by 8.

8	169	
8	21	1
	2	5

so,

$$(169)_{10} = (251)_8$$

Ans

i/ $2A7D_{16} = (?)_8$

First we will convert it to binary with the help of table.

<u>2</u>	<u>A</u>	<u>7</u>	<u>D</u>
0010	1010	0111	1101

now convert the binary numbers to octal by making groups of 3

<u>000</u>	<u>010</u>	<u>101</u>	<u>001</u>	<u>111</u>	<u>101</u>
0	2	5	1	7	5

so,

$$(2A7D)_{16} = (25175)_8$$

Ans/1/

$$J, \quad 11111111_2 = \pm (?)_{10}$$

Using weighted notation of magnitude bits.

$$\begin{aligned} &= 1 \times 2^6 + 1 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 \\ &= 64 + 32 + 16 + 8 + 4 + 2 + 1 \\ &= (127)_{10} \end{aligned}$$

So, the sign-bit is 1.

hence

$$(11111111)_2 = (-127)_{10}$$

Ans

$$K, \quad -12_{10} = (?)_2$$

Using repeated division.

2	12	
2	6	0
2	3	0
	1	1

$$(12)_{10} = (1100)_2 \quad \text{or} \quad (00001100)_2$$

Now taking 2's complement of obtained numbers

$$\begin{array}{r} 00001100 \\ 11110011 \\ + \quad \quad \quad 1 \\ \hline 11110100 \end{array}$$

1's complement
2's complement

30

$$(-12)_{10} = (11110100)_2$$

Ans

L/ $198_{10} = (?)_{BCD}$

Using Decimal - BCD table.

$$\begin{array}{r} 1 \quad 9 \quad 8 \\ \hline 0001 \quad 1001 \quad 1000 \end{array}$$

30

$$(198)_{10} = (000110011000)_{BCD}$$

Ans

n/ $10000110000_{BCD} = (?)_{10}$

Using BCD-Decimal table

<u>1000</u>	<u>0111</u>	<u>0000</u>
8	7	0

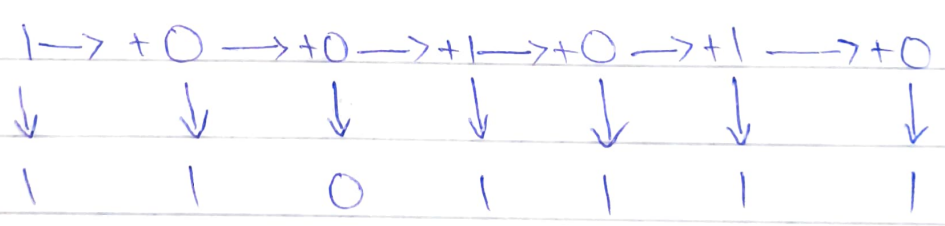
so

$(10000110000)_{BCD} = (870)_{10}$

Ans

n/ $(1001010)_2 = (?)_{Gray}$

Solution:

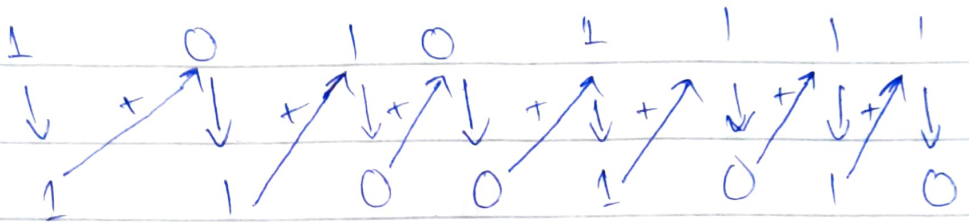


$(1001010)_2 = (1101111)_{Gray}$

Ans

Q, $10101111_{\text{Gray}} = (?)_2$

Sol:



80

$(10101111)_{\text{Gray}} = (11001010)_2$

(Ans)

P, $01000001 = (?)_{\text{ASCII}}$

Using ASCII table

$= 1 \times 2^6 + 1 \times 2^0$

$= 64 + 1$

$= 65_{\text{Hex}}$

~~$(0010000001)_{\text{Hex}} = 2100000001$~~

$(64)_{10} = (A)_{\text{ASCII}}$

(Ans)

Q, 111000 = (?111000)_{Even parity.}

since in an even parity there should be even amount of 1's, so we add 1 to the given number.

= (1111000)_{Even parity}

(Ans)

Q9, Calculate each of the following:

$$a) 0111111_2 - 00000111_2$$

Taking 2's complement on

$$\begin{array}{r} 00000111 \\ 11111000 \quad \text{1's complement} \\ + 1 \quad \text{2's complement} \\ \hline 11111001 \end{array}$$

Now,

$$\begin{array}{r} 10111111 \\ + 11111001 \\ \hline 101111000 \end{array}$$

discarded bit

so,

$$(01111000)_2$$

Answer
=

$$\text{Ex } 01101010_2 \times 11110001_2$$

Using 2's complement.

$$\begin{array}{r}
 \underline{11110001} \\
 00001110 \quad \text{1's complement} \\
 \hline
 1 \quad \text{2's complement} \\
 00001111
 \end{array}$$

Now,

$$\begin{array}{r}
 00001111 \\
 \underline{01101010} \\
 00000000 \\
 00001111 \quad x \\
 00000000 \quad x \quad x \\
 00001111 \quad x \quad x \quad x \\
 00000000 \quad x \quad x \quad x \quad x \\
 00001111 \quad x \quad x \quad x \quad x \quad x \\
 \underline{00000000 \quad x \quad x \quad x \quad x \quad x \quad x} \\
 000011000110 \quad 0110
 \end{array}$$

So,

$$(11000110110)_2$$

Ans

$$C_1 \quad 10001000_2 \div 001000010_2$$

$$\text{Quotient} = 00000000$$

Using 2's complement.

$$\begin{array}{r}
 00100010 \\
 11011101 \quad \text{1's complement} \\
 \hline
 1 \quad \text{2's complement} \\
 11011110
 \end{array}$$

Now,

$$\begin{array}{r}
 10001000 \\
 + 11011110 \\
 \hline
 10110010
 \end{array}$$

discarded bit \rightarrow

Adding 1 to the quotient.
so, 00000001

Subtracting divisor from 1st remainder using 2's complement.

$$\begin{array}{r}
 01100110 \\
 + 11011110 \\
 \hline
 101000100
 \end{array}$$

discarded \rightarrow

Adding 1 to the quotient = 00000010

Again

$$\begin{array}{r} \overset{1}{0} \overset{1}{0} \overset{1}{0} \overset{1}{0} \overset{1}{0} \overset{1}{0} \\ + \\ \hline 1101110 \end{array}$$

Discard \leftarrow 100100010

Adding 1 to quotient

$$= 00000011$$

New again

$$\begin{array}{r} + 00100010 \\ \hline 1101110 \end{array}$$

discard \leftarrow 10000000

Adding 1 to quotient

$$= 00000100$$

Answer.

Qr $6D_{16} - 3F_{16}$

First converting both hex numbers to binary numbers.

$$\begin{array}{cc} \underline{3} & \underline{F} \\ 0011 & 1111 \end{array}$$

now,

$$\begin{array}{cc} \underline{6} & \underline{D} \\ 0110 & 1001 \\ (6D)_{16} = (01101001)_2 \end{array}$$

Taking 9's complement of 3F.

$$\begin{array}{r} \underline{00111111} \\ 11000000 \quad \text{1's complement} \\ \hline 1 \quad \text{9's complement} \\ 11000001 \end{array}$$

$$\begin{array}{cc} \underline{1100} & \underline{0001} \\ e & 1 \end{array}$$

$$\begin{array}{r} 6D \\ + e1 \\ \hline 19E \end{array}$$

so

$$= 19E$$

ANSWER

$$e_1 \quad 00010110_{BCD} + 00010101_{BCD} = (?)_{10}$$

0001 0110

00 01 0101

0010 1010 invalid due to (>9).

Add 6 to the invalid code =

0010 1010

+ 0110

0011 0001

~~0011~~

0011

3

0001

1

(Any)

Q3, Apply CRC to the data bits
 11010011_2 using the generator code
 1010_2 to produce the transmitted
 CRC code.

$$D: 11010011$$

$$G: 1010$$

Since, the generator code has four data bits, add four 0's to the data byte. The appended data is

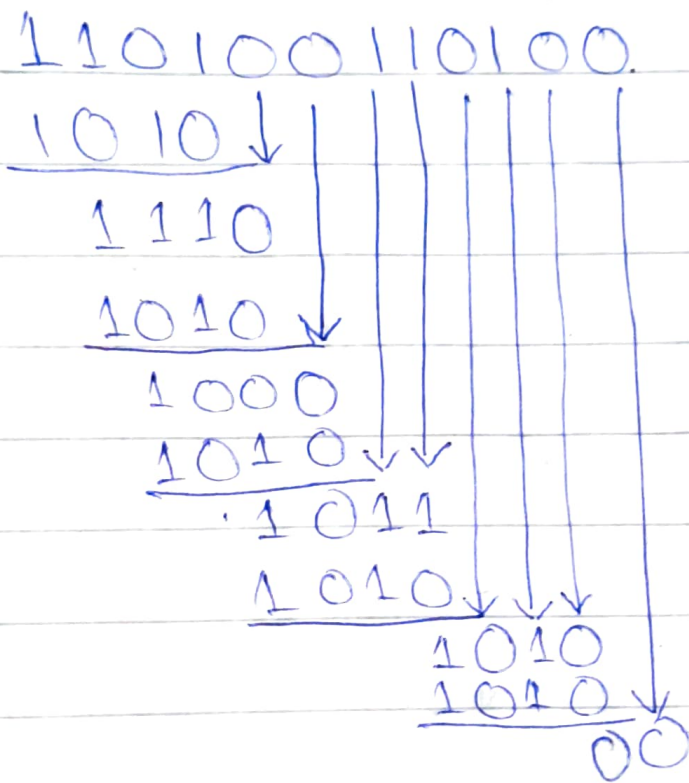
$$D' = 110100110000$$

Divide the appended data by the generator code using the modulo-2 operation until all bits are used.

$$\frac{D'}{G} = \frac{110100110000}{1010}$$

$$\begin{array}{r}
 110100110000 \\
 \underline{1010} \downarrow \\
 1110 \\
 \underline{1010} \downarrow \\
 1000 \\
 \underline{1010} \downarrow \\
 1011 \\
 \underline{1010} \downarrow \\
 1000 \\
 \underline{1010} \downarrow \\
 100
 \end{array}$$

Remainder = 0100. Since the remainder is not 0, append the data with the four remainder bits. Then divide by the generator code. The transmitted CRC is 110100110100.



Remainder = 0.

110100110100 is CRC code.

Ans

Q4 Assume that the code produced in problem Q.3 incurs an error in the most significant bit during transmission. Apply CRC to detect error.

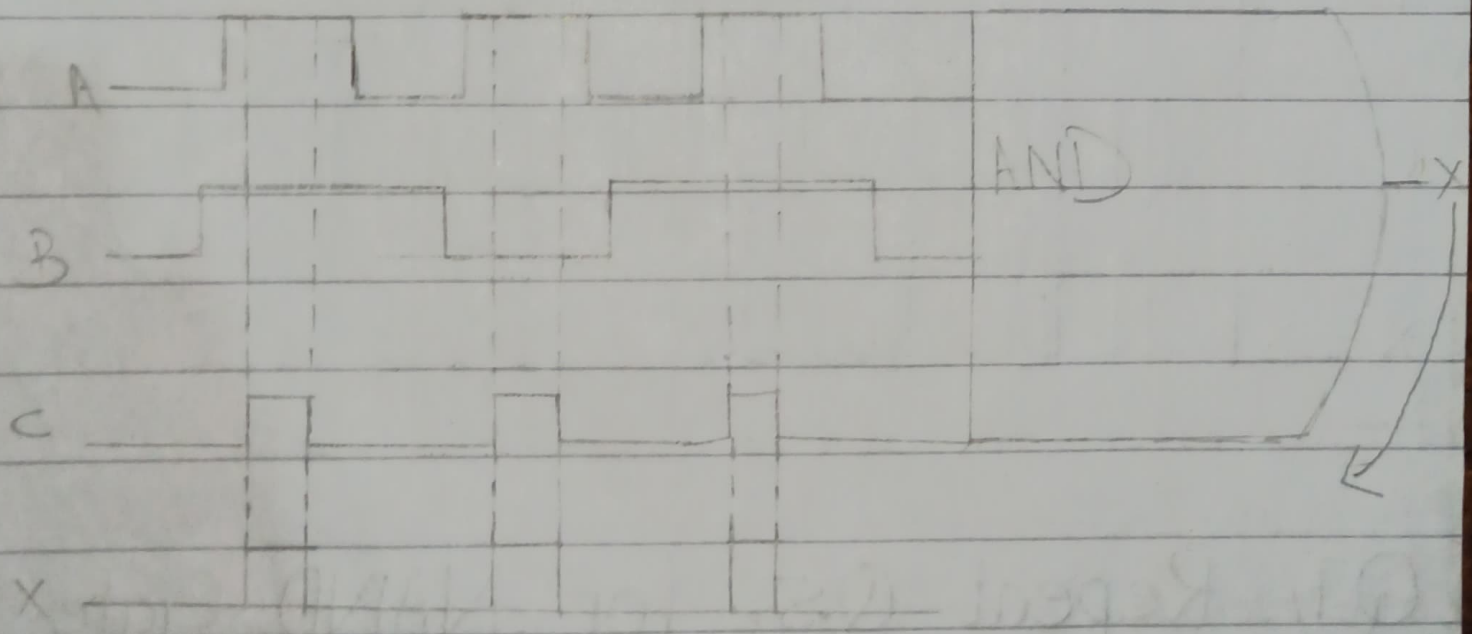
$$D = 010100110100$$

$$G = 1010$$

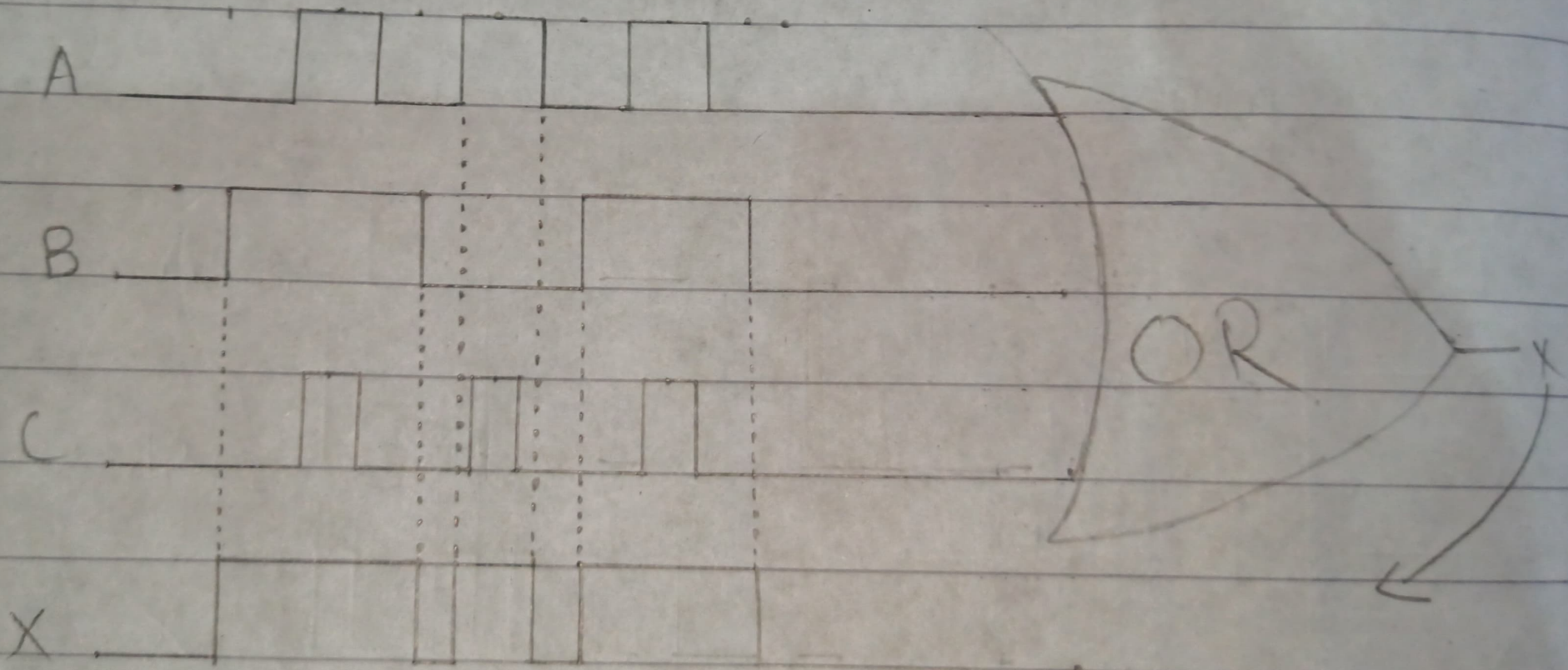
$$\begin{array}{r}
 010100110100 \\
 \underline{1010} \\
 1111 \\
 \underline{1010} \downarrow \\
 01010 \\
 \underline{1010} \downarrow \downarrow \downarrow \\
 000000110 \\
 \underline{1010} \\
 1100 \\
 \underline{1010} \downarrow \\
 01101 \\
 \underline{1010} \downarrow \\
 01110 \\
 \underline{1010} \downarrow \\
 01000 \\
 \underline{1010} \\
 0010
 \end{array}$$

Remainder $\neq 0$
 Hence, Error (~~is~~ ~~detected~~)
 is detected.

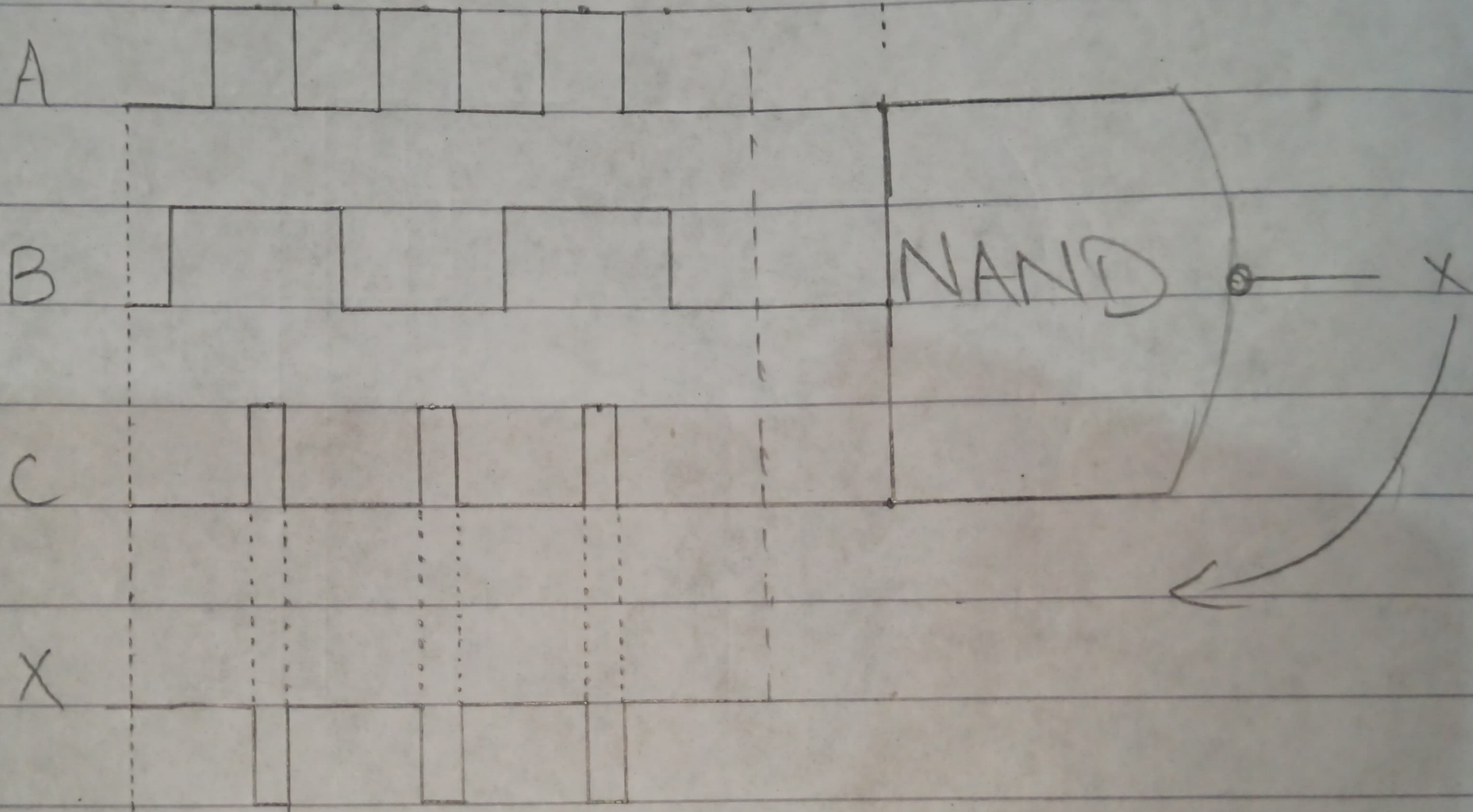
Q5. The input waveforms in Fig 1 is applied to a 3-input AND gate. Show the output waveform in proper relation to the inputs within a timing diagram.



Q6, Repeat Q.5 for a 3-input OR gate.

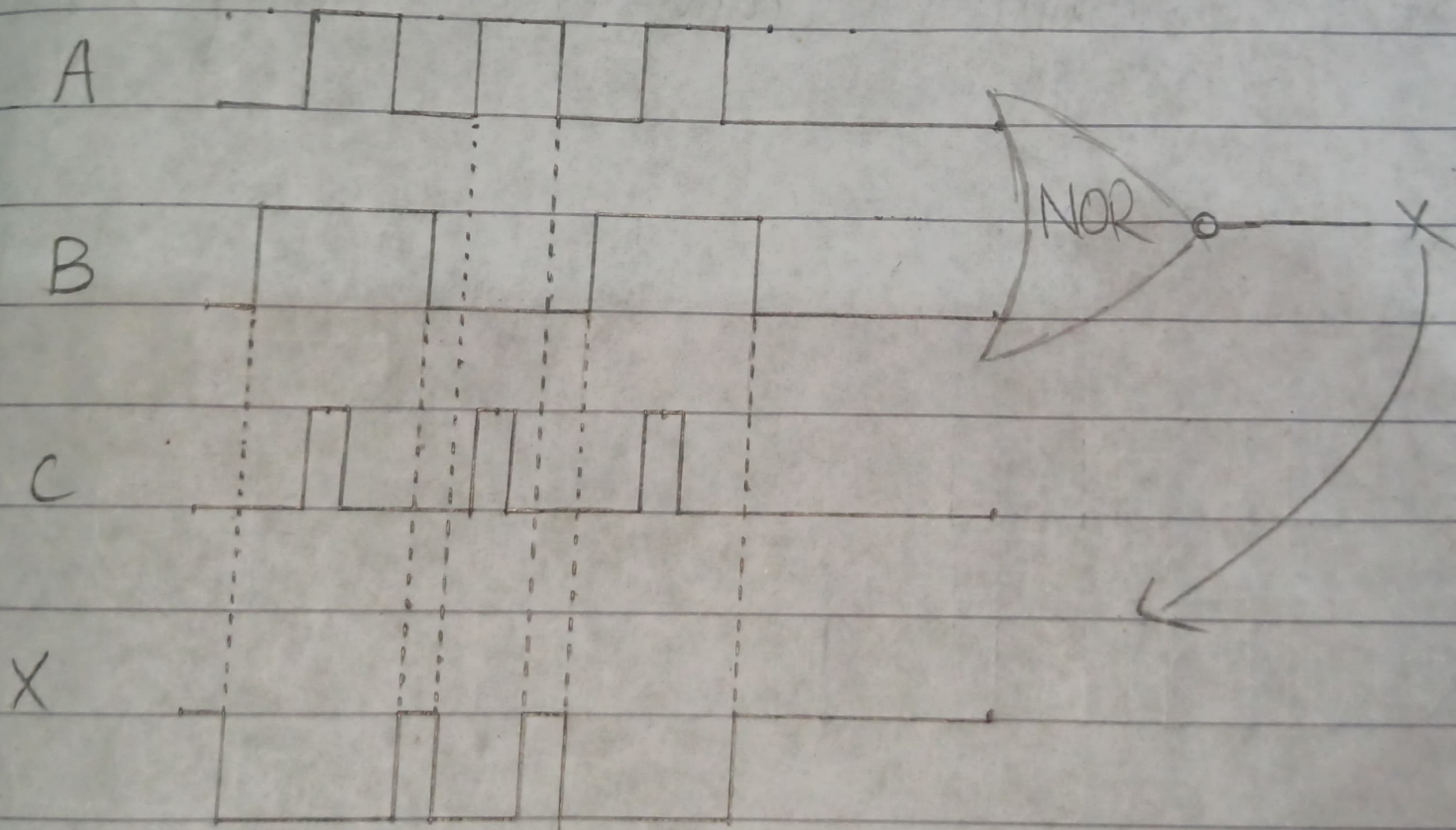


Q7, Repeat Q.5, for NAND Gate.

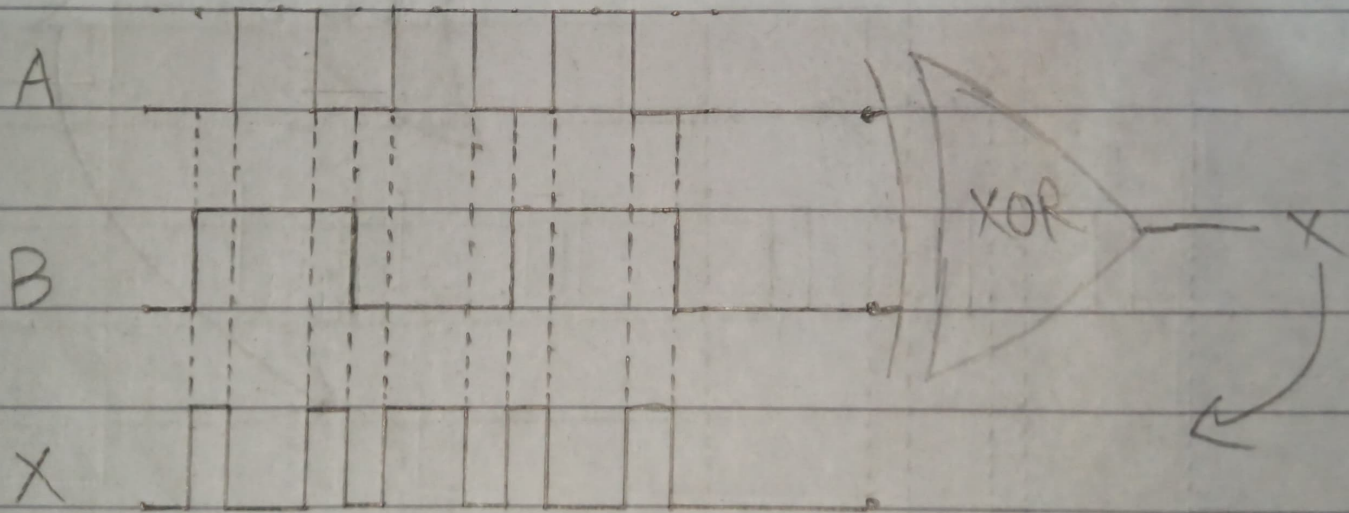


Q8: Repeat Gate.

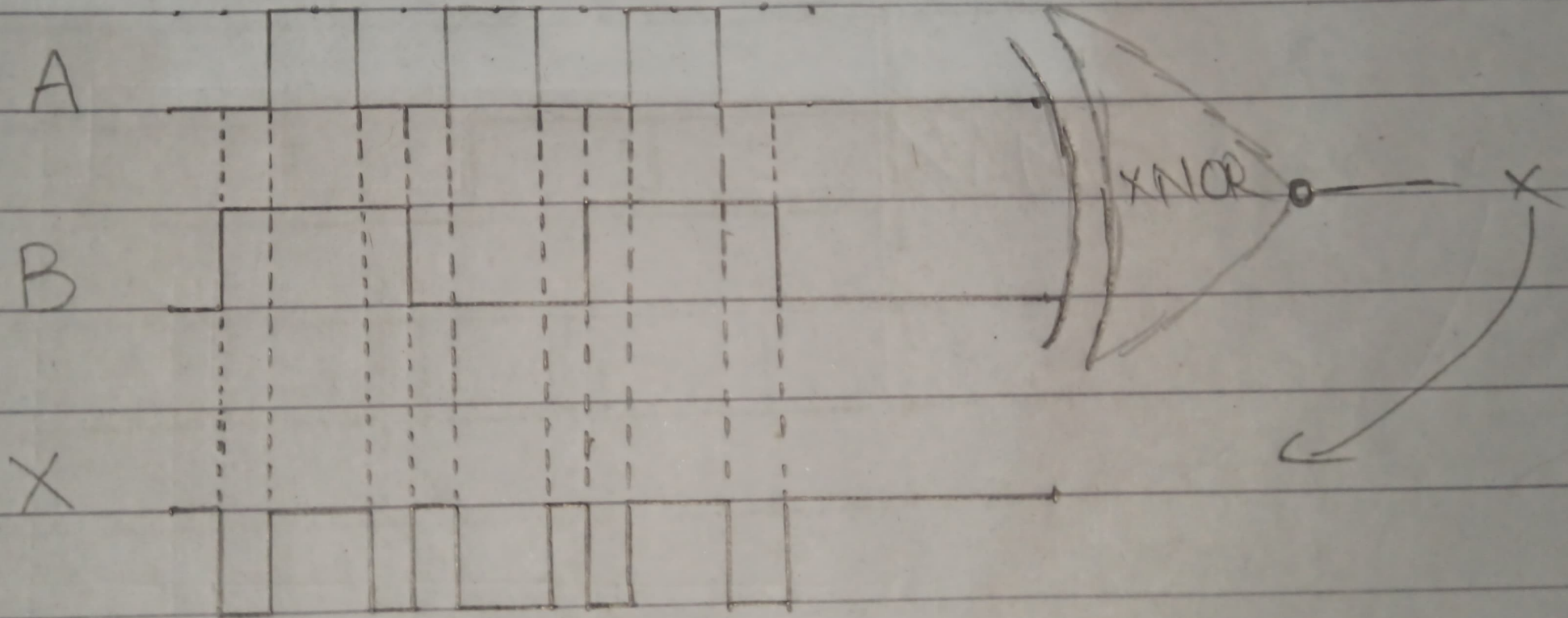
Q5, for NOR



Q9, The input waveforms in Fig 9 is applied to a XOR gate. Show the output waveform in proper relation to the inputs with a timing diagram.



Q10, Repeat Q.9 for XNOR gate.



Q11, Using Boolean algebra techniques, simplify the following expressions as much as possible.

$$\bar{A}\bar{B} + \bar{A}\bar{B}C + \bar{A}\bar{B}CD + \bar{A}\bar{B}CDE$$

Using Boolean Algebra Rules:

$$= \bar{A}\bar{B} + \bar{A}\bar{B}C + \bar{A}\bar{B}CD + \bar{A}\bar{B}CDE$$

$$= \bar{A}\bar{B} + \bar{A}\bar{B}CD + \bar{A}\bar{B}CDE$$

$$= \bar{A}\bar{B} + \bar{A}\bar{B}(CDE)$$

$$= \bar{A}\bar{B}$$

Ans

RAF WORK

$$\bar{A}\bar{B} + \bar{A}\bar{B}C$$

$$\bar{A}\bar{B}(1+C)$$

$$A+1=1$$

$$\bar{A}\bar{B}(1)$$

$$(\bar{A}\bar{B})$$

Q12. Convert the following expressions to standard SOP forms: $(C+D)(\bar{A}+D)$

First convert the given expression to SOP form

$$(C+D)(\bar{A}+D)$$

Distributing.

$$C\bar{A} + CD + D\bar{A} + DD$$

Domain of this SOP is ACD

Term $C\bar{A}$ is missing D

$$\Rightarrow C\bar{A} = C\bar{A}(D + \bar{D}) = C\bar{A}D + C\bar{A}\bar{D}$$

Term CD is missing A

$$\Rightarrow CD = CD(A + \bar{A}) = CDA + CD\bar{A}$$

Term $D\bar{A}$ is missing C

$$\Rightarrow D\bar{A} = D\bar{A}(C + \bar{C}) = D\bar{A}C + D\bar{A}\bar{C}$$

Term \overline{DD} is missing A & C.

$$\Rightarrow \overline{DD} = \overline{DD}(A + \overline{A}) = \overline{DD}A + \overline{DD}\overline{A}$$

Term $\overline{DD}A$ & $\overline{DD}\overline{A}$ is missing C

$$\Rightarrow \overline{DD}A = \overline{DD}A(C + \overline{C}) = \overline{DD}AC + \overline{DD}A\overline{C}$$

$$\Rightarrow \overline{DD}\overline{A} = \overline{DD}\overline{A}(C + \overline{C}) = \overline{DD}\overline{A}C + \overline{DD}\overline{A}\overline{C}$$

hence resulting SOP form is:

~~$(C\overline{A}\overline{D} + C\overline{A}D + CDA + CD\overline{A} + \overline{D}\overline{A}C + \overline{D}\overline{A}\overline{C} + \overline{D}D\overline{A}C + \overline{D}D\overline{A}\overline{C} + \overline{D}D\overline{A}C + \overline{D}D\overline{A}\overline{C})$~~

$$(C\overline{A}\overline{D} + C\overline{A}D + CDA + CD\overline{A} + \overline{D}\overline{A}C + \overline{D}\overline{A}\overline{C} + \overline{D}D\overline{A}C + \overline{D}D\overline{A}\overline{C} + \overline{D}D\overline{A}C + \overline{D}D\overline{A}\overline{C})$$

(Ans)

Q3. Write the standard POS expression using the standard ~~the~~ SOP expression obtained in Q.12.

$$\bar{C}\bar{A}D + \bar{C}\bar{A}\bar{D} + D\bar{A}\bar{C} + ACD.$$

Evolution of the POS expression is

$$(101) + (100) + (110) + (111)$$

In the domain of this expression there are three variables, there are $2^3 = 8$ possible combinations. Four of which are contained by this expression the rest are:

$$000, 010, 011, 001$$

hence, the equivalent POS expression is

$$(A+C+D)(A+\bar{C}+D)(A+\bar{C}+\bar{D})(A+C+\bar{D})$$

Q14/ Draw a ~~single~~ single truth table for both the standard POS and SOP expression obtained in Q.12 & Q.13.

A	C	D	X	POS/SOP
0	0	0	0	$(A+C+D)$
0	0	1	0	$(A+C+\bar{D})$
0	1	0	0	$(A+\bar{C}+D)$
0	1	1	0	$(A+\bar{C}+\bar{D})$
1	0	0	1	$(A\bar{C}\bar{D})$
1	0	1	1	$(A\bar{C}D)$
1	1	0	1	$(AC\bar{D})$
1	1	1	1	(ACD)

POS Expression:

$$(A+C+D)(A+\bar{C}+D)(A+\bar{C}+\bar{D})(A+C+\bar{D})$$

SOP Expression:

$$(\bar{A}\bar{D}) + (\bar{C}\bar{A}\bar{D}) + (CDA) + (DA\bar{C})$$

Q15. Use a karnaugh map to simplify the following expressions to a minimum SOP form:

$$\bar{A}\bar{B}\bar{C} + \bar{A}BC + A\bar{B}\bar{C} + ABC$$

$$000 + 011 + 101 + 110$$

AB \ C	0	1	
00	1		$\rightarrow (\bar{A}\bar{B}\bar{C})$
01		1	$\rightarrow (\bar{A}BC)$
11	1		$\rightarrow (A\bar{B}\bar{C})$
10		1	$\rightarrow (A\bar{B}C)$

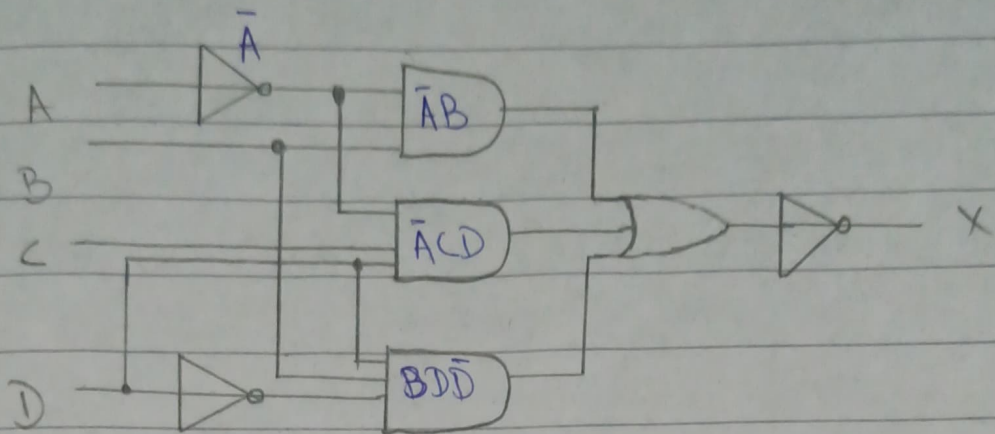
$(\bar{A}\bar{B}\bar{C}) + (\bar{A}BC) + (A\bar{B}\bar{C}) + (A\bar{B}C)$ is the minimum SOP.

Q16. Obtain the minimum POS expression from the karnaugh map used in Q.15.

AB \ C	0	1	
00	1	0	$\rightarrow (A+B+\bar{C})$
01	0	1	$\rightarrow (A+\bar{B}+C)$
11	1	0	$\rightarrow (A+B+\bar{C})$
10	0	1	$\rightarrow (\bar{A}+B+C)$

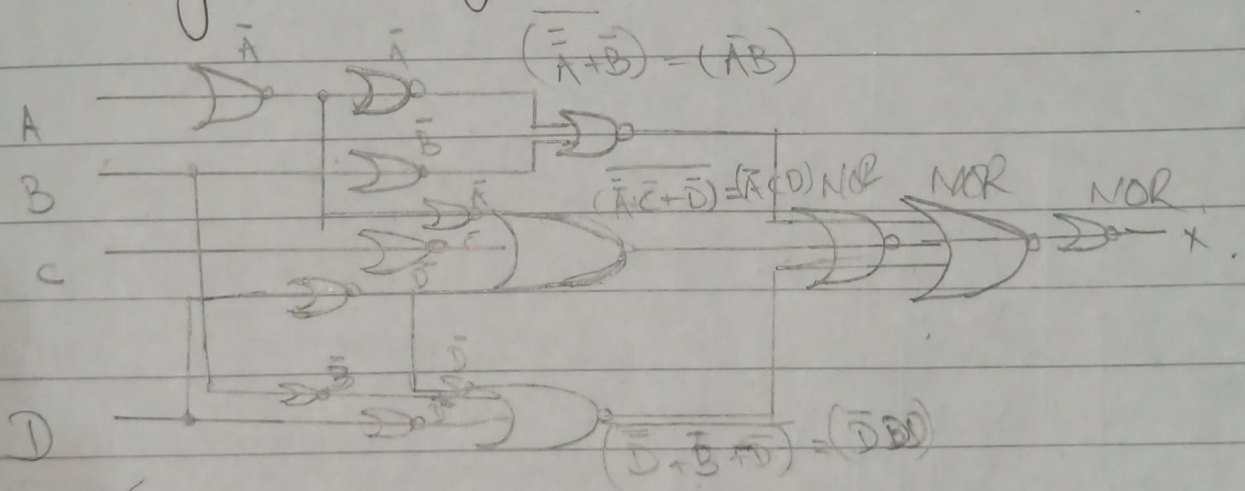
$(A+B+\bar{C})(A+\bar{B}+C)(A+B+\bar{C})(\bar{A}+B+C)$ is the minimum POS expression.

Q17, Write the output expression for circuit in Figure 3.



$$X = (\bar{A}B) + (\bar{A}CD) + (BD\bar{D})$$

Q18, Implement the logic circuits in Fig 3 using only NOR gates.



$$X = (\bar{A}+B) \cdot (\bar{A}+\bar{C}+\bar{D}) \cdot (\bar{D}+\bar{B}+\bar{D}) = (\bar{A}B) + (\bar{A}CD) + (\bar{D}BD)$$

Ans

Q.19 is also the same X. Solved on the last page.

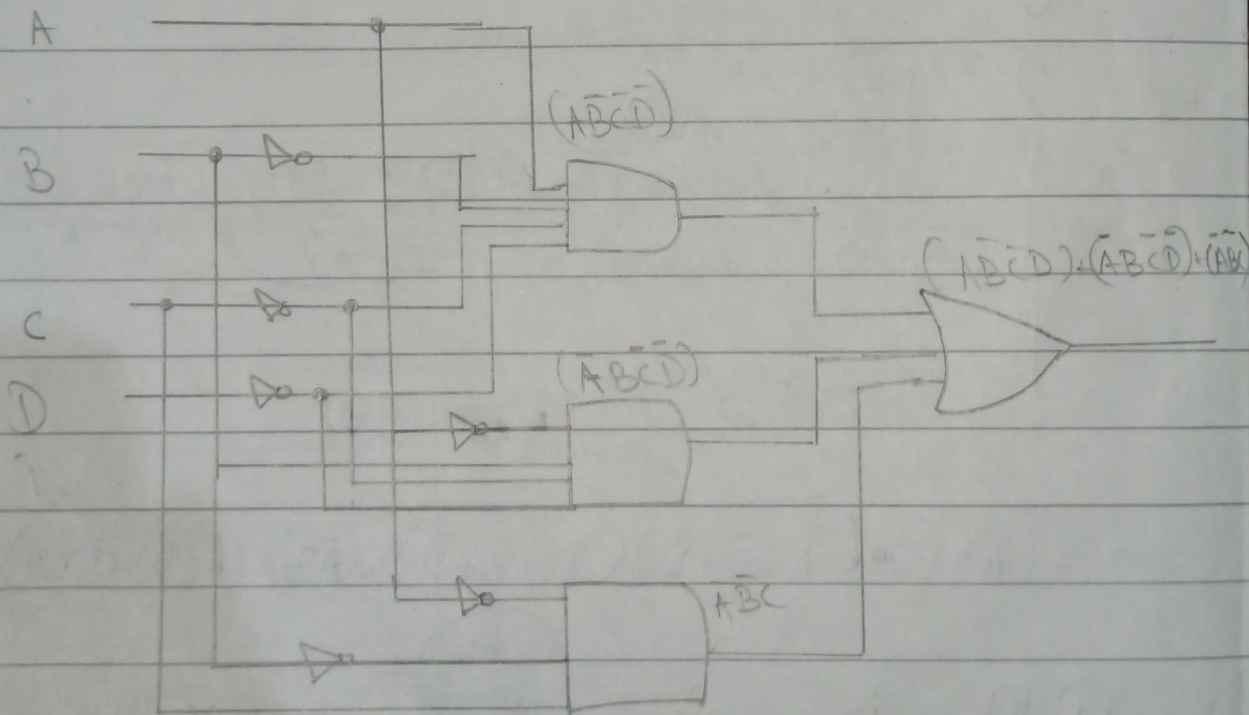
Q20, Implement a logic circuit for the truth table in Table 1.

obtained expression from the truth table is

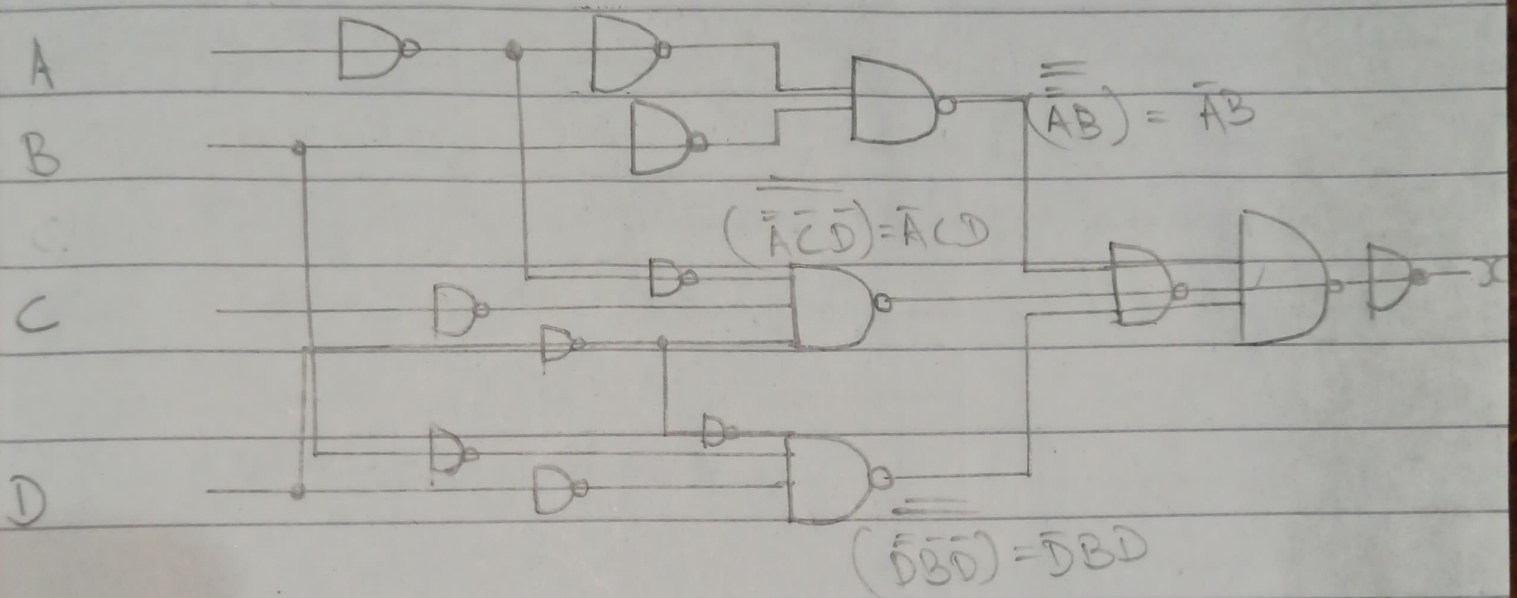
$$(\bar{A}\bar{B}\bar{C}\bar{D}) + (\bar{A}\bar{B}CD) + (\bar{A}B\bar{C}\bar{D}) + (A\bar{B}\bar{C}\bar{D}) + (A\bar{B}C\bar{D}) + (A\bar{B}C\bar{D}) + (A\bar{B}CD) + (ABCD)$$

By reducing the expression using boolean laws and rules we get

$$(\bar{A}\bar{B}\bar{C}\bar{D}) + (\bar{A}B\bar{C}\bar{D}) + (\bar{A}\bar{B}C)$$



Q19, Implement the logic circuits in Fig 3 using only NAND gates.



$$X = (\overline{AB}) + (\overline{\overline{A}CD}) + (\overline{\overline{D}BD})$$