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Id NO:- 17011

Class timming:- wednesday (8:am to 11:-00am)

Subject:- Discrete mathematics

Program:- BS(CS)

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"Question no 1"

what is Recurrence Relation? And explain Repeated Substitution method with help of example.

Answer:

Recurrence Relation:

A recurrence relation is an equation that recursively defines a sequence where the next term is a function of the previous terms (Expressing  $F_n$  as some combination of  $F_i$  with  $i < n$ ). Example - Fibonacci Series  
 $F_n = F_{n-1} + F_{n-2}$

Repeated Substitution method:-

Example:-

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ 4T(n/2) + n & \text{if } n \geq 2 \end{cases} \quad n = 2^k$$

(a)  $T(n) = n + 1 \cdot (n/2)$

Let's write  $T(n)$  as a function of  $(n/2)$

$$T(n) = \frac{n}{2} + 4T\left(\frac{n}{2}\right) = \frac{n}{2} + 4T\left(\frac{n}{4}\right)$$

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$$T(n) = n + 4 \left[ \frac{n}{2} + 4T\left(\frac{n}{4}\right) \right] = n + 4 \cdot \frac{n}{2} + (4^2)T\left(\frac{n}{4}\right)$$

$$T(n) = n + 2n + 16T\left(\frac{n}{4}\right)$$

Let's write  $T(n)$  as a function of  $\left(\frac{n}{4}\right)$

$K=2$ ;

$$T(n) = \frac{n}{4} + 4T\left(\frac{n}{4}\right) = \frac{n}{4} + 4T\left(\frac{n}{8}\right)$$

$$T(n) = n + 2n + 16 \left[ \frac{n}{4} + 4T\left(\frac{n}{8}\right) \right]$$

$$= n + 2n + \frac{16n}{4} + (16)(4)T\left(\frac{n}{8}\right)$$

$$T(n) = n + 2n + 4n + 64T\left(\frac{n}{8}\right)$$

Let's write  $T(n)$  as a function of  $\left(\frac{n}{8}\right)$

$K=3$ ;

function of  $\left(\frac{n}{8}\right)$

$$T(n) = \frac{n}{4} + 4T\left(\frac{n}{4}\right) = \frac{n}{4} + 4T\left(\frac{n}{8}\right)$$

$$T(n) = n + 2n + 4n + 64T\left(\frac{n}{8}\right)$$

Let's write  $T(n)$  as a function of  $\left(\frac{n}{8}\right)$

$$T(n) = \frac{n}{8} + 4T\left(\frac{n}{8}\right) = \frac{n}{8} + 4T\left(\frac{n}{16}\right)$$

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$$T(n) = n + 2n + 4n + 64 \left[ \frac{n}{8} + 4T(n/16) \right]$$
$$= n + 2n + 4n + \frac{64n}{8} + (64)(4)T(n/16)$$

$$T(n) = n + 2n + 4n + 8n + 256T(n/16)$$

At this point we can see a pattern so we stop "repeating substitution" and write a general expression that describes the behavior.

$$(a) T(n) = n + 2n + 4n + 8n + 256T(n/16)$$
$$T(n) = n(1 + 2 + 4 + 8) + 256T(n/16)$$
$$T(n) = n(2^0 + 2^1 + 2^2 + 2^3) + (4)^4 T(n/2^4)$$

we can now generalize as follows

$$T(n) = n(2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{k-1}) + 4^k T(n/2^k)$$
$$T(n) = n \sum_{i=0}^{k-1} 2^i + 4^k T(n/2^k)$$

The problem says that  $n = 2^k$  so,

$$T(n) = n \sum_{i=0}^{k-1} 2^i + 4^k T(2^k/2^k)$$

$$T(n) = n \sum_{i=0}^{k-1} 2^i + 4^k T(1)$$

The problem says that  $T(n) = 1$  if  $n = 1$

$$T(n) = n \sum_{i=0}^{k-1} 2^i + 4^k (1)$$

$$T(n) = n \sum_{i=0}^{k-1} 2^i + 4^k$$

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## "Question no 2"

### Premises:

If there was a ball game, then traveling was difficult. If they arrived on time, then traveling was not difficult. They arrived on time.

### Conclusion:-

There was no ball game.

Determine symbolically using rule of Inference.

P: There was a ball game.  
Q: Traveling was difficult.  
R: They arrived on time.

### Solution:-

|     |                             |                   |
|-----|-----------------------------|-------------------|
| 1 → | $P \rightarrow Q$           | Premises          |
| 2 → | $R \rightarrow \sim Q$      | Premises          |
| 3 → | R                           | Premises          |
| 4 → | $\sim Q$                    | 2, 3 modus ponens |
| 5 → | $\sim Q \rightarrow \sim P$ | 1, Contrapositive |

$\sim P$

4, 5 Modus ponens

### Conclusion:-

$\sim P$  = There was no ball game.

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"Question no 3"

Determine logically Symbolically using  
rule of inference.

Solution:-

P: Claghorn has wide support.

Q: Claghorn is asked to run for the Senate

X: Claghorn yells "Eureka"

$P \rightarrow Q$

Premises

$X \rightarrow \sim Q$

Premises

X

Premises

$\sim Q$

2,3 Modus ponens

$\sim Q \rightarrow \sim P$

1, Contrapositive

---

$\sim P$

---

5,6 modus ponens

Conclusion:-

$\sim P$  = Claghorn does not have  
wide support.

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"Question no 4"

Answer:-

Pigeon hole principle:-

The pigeon hole principle states, that if  $n$  items are put into  $m$  containers, with  $n > m$ , then at least one container must contain more than one item. If you have more "objects" than you have "holes", at least one hole must have multiple objects in it. A real-life example could be "if you have three gloves, then you have at least two right hand gloves, or at least two left hand gloves, because you have three objects, but only two categories to put them into (right to left). The seemingly obvious statement, a type of counting argument, can be used to demonstrate possibly unexpected results. For example if you know that the population of London is greater by (at least two people) than the maximum number of hairs

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than can be present on a human head, then the pigeon hole principle requires that there must be (at least) two people in London who have the same number of hairs on their heads.

The principle has several generalizations and can be stated in various ways.

In a more quantified version. For natural numbers  $k$  and  $m$ , if  $n = km + 1$  objects are distributed among  $m$  sets. then the pigeonhole principle asserts that at least one of the sets will contain at least  $k+1$  objects. For arbitrary  $n$  and  $m$  generalizes to  $k+1 = \lfloor (n-1)/m \rfloor + 1 = \lfloor n/m \rfloor + 1$ .



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### Example for pigeonhole:-

- 1) Finding maximum and minimum value in array. Let the minimum and maximum values be "min" and "max" respectively.
- 2) Also find range as " $\text{max} - \text{min} + 1$ ".
- 3) visit each element of the array then put each element in its pigeonhole.
- 4) Start the loop all over the pigeonhole.

Input data:

8 | 3 | 2 | 7 | 4 | 6 | 8

|   |  |  |  |  |  |  |  |                              |
|---|--|--|--|--|--|--|--|------------------------------|
| 0 |  |  |  |  |  |  |  | Range: $8 - 2 + 1 = 7$       |
| 1 |  |  |  |  |  |  |  | $\text{arr}[i] - \text{min}$ |
| 2 |  |  |  |  |  |  |  |                              |
| 3 |  |  |  |  |  |  |  |                              |
| 4 |  |  |  |  |  |  |  | $8 - 2 = 6$                  |
| 5 |  |  |  |  |  |  |  |                              |
| 6 |  |  |  |  |  |  |  |                              |

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|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 8 | 3 | 2 | 7 | 4 | 6 | 8 |
|---|---|---|---|---|---|---|

|   |     |
|---|-----|
| 0 | 2   |
| 1 | 3   |
| 2 | 4   |
| 3 |     |
| 4 | 6   |
| 5 | 7   |
| 6 | 8 8 |

Range:  $8 - 2 + 1 = 7$

$arr[i] - \min$

$3 - 2 = 1$

Now

$2 - 2 = 0$

Now:-

Subtract third away

$arr[i] - \min$

$7 - 2 = 5$

Now:-

$arr[i] - \min$

$4 - 2 = 2$

Now:-

$arr[i] - \min$

$6 - 2 = 4$

Now:-

$8 - 2 = 6$

Now this is the example of region hole the in one hole there are two regions init.

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"Question no 6"

Answer:-

Let  $P$ ,  $Q$  and  $R$  be the three Control panels.

Let  $S$  be the System.

Let 0 represent that the Switch in the Control panel is in the off ~~possib~~ position.

Let 1 represents that the System is disabled and  $S=1$  represents that the System is enabled.

If two or three of the Control panels have the Switches in the on position then the System is enabled.

$P=0$  and  $Q=1$  and  $R=1$

$P=1$  and  $Q=0$  and  $R=1$ ,  
then  $S=1$ .

If one or ~~more~~ none of the Control panels have the Switches in the on ~~Switches~~ position, then the System is disabled.

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$p=1$  and  $Q=0$  and  $R=0$ , then  $S=0$   
 $p=0$  and  $Q=1$  and  $R=0$ , then  $S=0$   
 $p=0$  and  $Q=0$  and  $R=1$ , then  $S=0$   
 $p=0$  and  $Q=0$  and  $R=0$ , then  $S=0$

Combining this information in an input/output table, we then obtain:

| Input |   |   | output |
|-------|---|---|--------|
| P     | Q | R | S      |
| 1     | 1 | 1 | 1      |
| 1     | 1 | 0 | 1      |
| 1     | 0 | 1 | 1      |
| 1     | 0 | 0 | 0      |
| 0     | 1 | 1 | 1      |
| 0     | 1 | 0 | 0      |
| 0     | 0 | 1 | 0      |
| 0     | 0 | 0 | 0      |

Boolean expression:-

Negation  $\sim p$ : not p

Disjunction  $p \vee q$ :  $p$  or  $q$

Conjunction  $p \wedge q$ : p and q

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If two or the three Control panels have the switches in the on positions ( $S=1$ ) then their  $P \wedge Q$ ,  $P \wedge R$ , or  $Q \wedge R$  is equal to 1 (or all three are true at the same time when all three Control panels have the switches in the on position). The boolean expression is then the disjunction of these three statements

$$(P \wedge Q) \vee (Q \wedge R) \vee (P \wedge R)$$

Circuit:-

A "not gate" outputs the negation  $\sim a$  of the input  $a$ .

An "OR Gate" outputs the disjunction  $a \vee b$  of the inputs  $a$  and  $b$ .

An AND Gate outputs the conjunction  $a \wedge b$  of the inputs  $a$  and  $b$ .

$$(P \wedge \sim Q) \vee (\sim P \wedge Q)$$

The negation will occur before the conjunctions and disjunction, thus we first negate  $P$  by using a NOT gate with input  $P$  and we negate

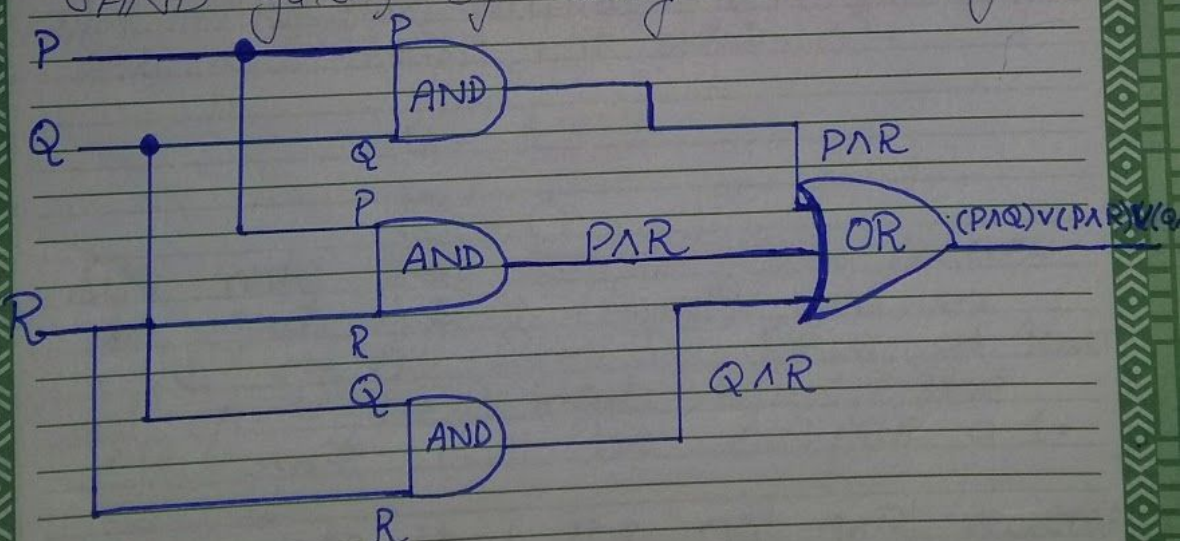
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Q by using a NOT gate with input

Next the conjunction will occur before the disjunction in  $(P \wedge \sim Q) \vee (\sim P \wedge Q)$  due to the two pairs of baskets. we then need to use an AND gate with input P and  $\sim Q$  to form  $P \wedge \sim Q$  while we use an AND gate with input  $\sim P$  and Q to form  $\sim P \wedge Q$ .

Finally, we need to take the disjunction of  $P \wedge \sim Q$  and  $\sim P \wedge Q$  (output AND gates) by using an OR gate.



Result:-

$$(P \wedge \sim Q) \vee (\sim P \wedge Q) \vee R$$

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"Question no 7"

Give the output signals for the circuit.

1) Answer:-

A Not gate changes the input from a 0 to a 1 or from a 1 to a 0.

An OR gate outputs a 1 if at least one of the inputs is a 1 (else it outputs 0).

An AND gate outputs a 1 if both inputs are a 1 (else it outputs 0).

$$P = 1$$

$$Q = 0$$

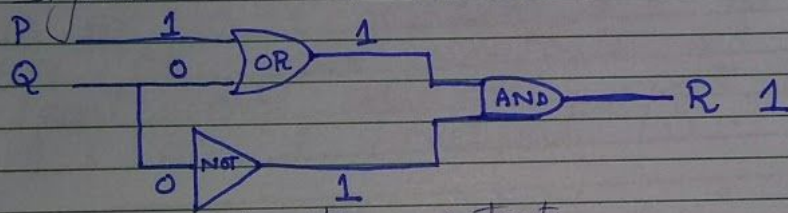
We note that the OR gate receives an input of 1 (from P) and an input of 0 (from Q). Since at least one of the inputs is a 1, the OR gate will be output a 1.

We note that the "not gate" receives an input of 0 (from Q), thus the NOT gate will output a 1.

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Finally, we then note that the AND gate receives an input of 1 from the OR gate and an input 1 from the NOT gate. As both inputs are a 1, the AND gate will output a 1 as well thus the output signal is a  $R=1$ .



| Inputs |   | Outputs |
|--------|---|---------|
| P      | Q | R       |
| 1      | 0 | 0       |
| 1      | 1 | 1       |
| 0      | 0 | 0       |
| 0      | 1 | 0       |



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Part II:-

input signals  $p=1$ ,  $Q=0$ ,  $R=0$

Answer:-

A "Not gate" changes the input from a 0 to a 1 or from a 1 to a 0.  
An "OR gate" outputs a 1 if at least one of the input is a 1 (else it outputs 0).

A "AND gate" outputs a 1 if both inputs are a 1 (else it outputs a 0).

$$p = 1, Q = 0, R = 0$$

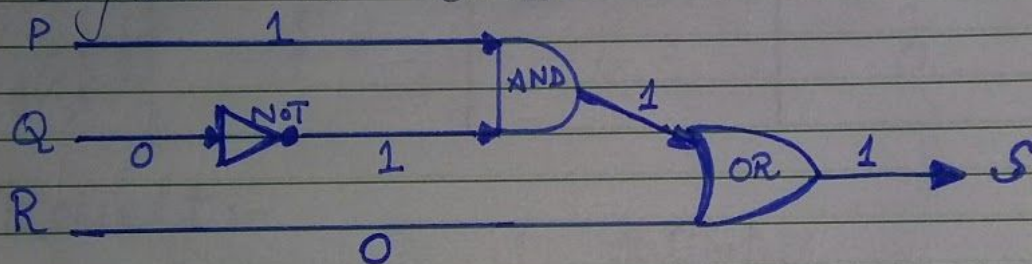
we note that NOT gate receives an input 0 (from  $Q$ ). Thus the not gate is output a 1.

we note the AND gate receives an input of 1 (from  $p$ ) and an input of 1 (from  $Q$ ). Since both inputs are a 1, the AND gate will output a 1.

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Finally, we note that the OR gate receives an input of 1 from the AND gate and an input of 0 from R. As at least one of the inputs is a 1, the OR gate will output a 1 as well thus the output signal is  $S=1$ .



Similarly, we can derive the output of each gate for every possible input P, Q and R.

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| P | Q | R | output NOT Gate | output AND Gate | output OR Gate |
|---|---|---|-----------------|-----------------|----------------|
| 1 | 1 | 1 | 0               | 0               | 1              |
| 1 | 1 | 0 | 0               | 0               | 0              |
| 1 | 0 | 1 | 1               | 1               | 1              |
| 1 | 0 | 0 | 1               | 1               | 1              |
| 0 | 1 | 1 | 0               | 0               | 1              |
| 0 | 1 | 0 | 0               | 0               | 0              |
| 0 | 0 | 1 | 1               | 0               | 1              |
| 0 | 0 | 0 | 1               | 0               | 0              |

Result:-

| P | Q | R | output S |
|---|---|---|----------|
| 1 | 1 | 1 | 1        |
| 1 | 1 | 0 | 0        |
| 1 | 0 | 1 | 1        |
| 1 | 0 | 0 | 1        |
| 0 | 1 | 1 | 1        |
| 0 | 1 | 0 | 0        |
| 0 | 0 | 1 | 1        |
| 0 | 0 | 0 | 0        |

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"Part 3"

input signals  $P=0$ ,  $Q=0$ ,  $R=0$

Answer:-

A "NOT gate" changes the input from a 0 to a 1 or from a 1 to 0.

An "OR gate" outputs a 1 if at least one of the inputs is a 1 (else it outputs 0).

An "AND gate" outputs a 1 if both inputs are a 1 (else it outputs 0).

we note that the leftmost OR gate receives an input of 0 (from P) and an input of 0 (from Q). Since at least one of the inputs is a 0, the OR gate will output a 0.

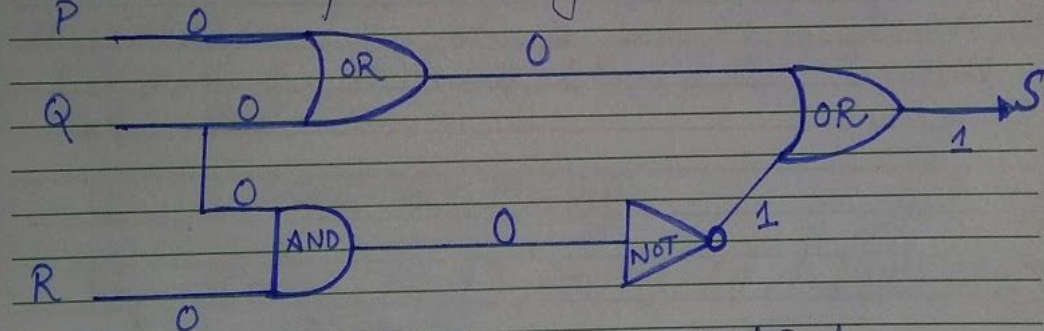
we note that the AND gate receives an input of 0 (from Q) and an input of 0 (from R). Since both inputs are a 0, the AND gate will output a 0.

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we then note that the NOT gate receives an input of 0 (from the AND gate), thus the NOT gate will output a 1.

Finally, we then note that the right most OR gate receives an input of 0 from the left most OR gate an input of 1 from the NOT gate. As at least one of the input is a 1, thus the OR gate will output a 1 as well and thus the output signal is  $S=1$ .



Result:-

| Input |   |   | Output |
|-------|---|---|--------|
| P     | Q | R | S      |
| 1     | 1 | 1 | 1      |
| 1     | 1 | 0 | 1      |
| 1     | 0 | 1 | 1      |
| 1     | 0 | 0 | 1      |
| 0     | 1 | 1 | 1      |
| 0     | 1 | 0 | 1      |
| 0     | 0 | 1 | 1      |
| 0     | 0 | 0 | 1      |

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"Question no 8"

Answer:-

$$A = \{0, 1, 2, 3\}$$

A relation  $R$  on a Set  $A$  is reflexive if  $(a, a) \in R$  for every element  $a \in A$ .

A relation  $R$  on a Set  $A$  is symmetric if  $(b, a) \in R$  whenever  $(a, b) \in R$ .

A relation  $R$  on a Set  $A$  is transitive if  $(a, b) \in R$  and  $(b, c) \in R$  implies  $(a, c) \in R$ .

$$A = \{0, 1, 2, 3\}$$

$$R_1 = \{(0, 0), (0, 1), (0, 3), (1, 1), (1, 0), (2, 3), (3, 3)\}$$

(a) Directed Graph:-

we note that  $A$  contains 4 elements and thus we will draw 4 points.

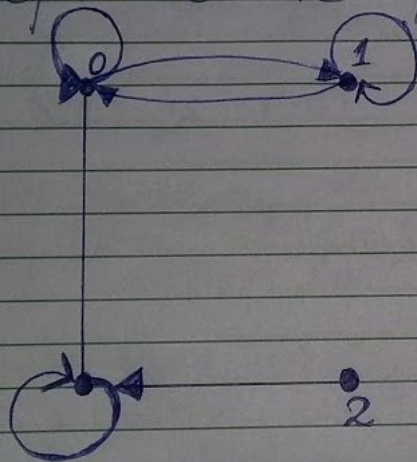
we label these points 0, 1, 2, 3 (which are the elements of  $A$ ).

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For every element  $(x, y) \in R$ , with  $x \neq y$ , we draw an arrow from  $x$  to  $y$ .

For every element  $(x, x) \in R$ , we draw a loop at the point  $x$ .



B) The relation  $R_1$  is reflexive if  $(a, a) \in R_1$  for every element  $a \in A$ .

Since  $A = \{0, 1, 2, 3\}$ ,  $R_1$  is reflexive if it contains  $(0, 0)$ ,  $(1, 1)$ ,  $(2, 2)$ ,  $(3, 3)$ .

we note that  $R_1$  does not contain  $(2, 2)$  and thus  $R_1$  is not reflexive.

" $R_1$  is not reflexive"

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C) The relation  $R_1$  on a set  $A$  is symmetric if  $(b, a) \in R_1$  whenever  $(a, b) \in R_1$

we note that  $(0, 3) \in R_1$ , while  $(3, 0) \notin R_1$  thus  $R_1$  is not symmetric.

" $R_1$  is not Reflexive"

D) The relation  $R_1$  on a set is transitive if  $(a, b) \in R_1$  and  $(b, c) \in R_1$  implies  $(a, c) \in R_1$

we note that  $(1, 0) \in R_1$  and  $(0, 3) \in R_1$ , while  $(1, 3) \notin R_1$  and thus  $R_1$  is not transitive.

" $R_1$  is not transitive"



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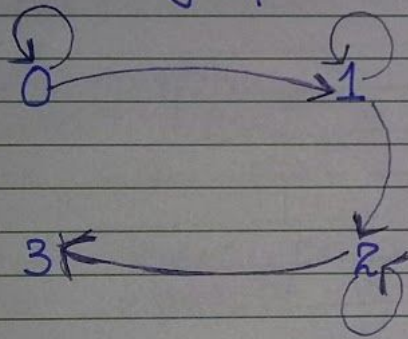
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$$R_2 = \{(0,0)(0,1)(1,1)(1,2)(2,2)(2,3)\}$$

$$A = \{0, 1, 2, 3\}$$

$$R_2 = \{(0,0)(0,1)(1,1)(1,2)(2,2)(2,3)\}$$

a) Directed graph:-



b) No,  $R_2$  is not reflexive because there is no loop at 3.

c) No,  $R_2$  is not symmetric because between two numbers where there exists an arrow, the arrow goes only one direction.

d) No,  $R_2$  is not transitive though there is an arrow between 0 and 1 and also an arrow between 1 and 2, but there is no arrow between 0 and 2.

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"Part 3"

$$R_3 = \{(2,3), (3,2)\}$$

Answer:-

$$A = \{0,1,2,3,4\}$$

$$R_3 = \{(2,3), (3,4)\}$$

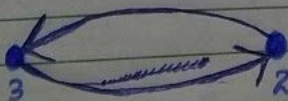
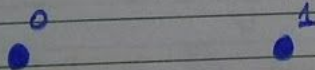
(a) Directed graph:-

we note that  $A$  contains 4 elements and thus we will draw 4 points

we label these points 0, 1, 2, 3 (which are the elements of  $A$ .)

For every element  $(x,y) \in R$  with  $x \neq y$  we draw an arrow from  $x$  to  $y$

For every element  $(x,x) \in R$ , we draw a loop at the point  $x$ .



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b) The relation  $R_3$  is reflexive if  $(a, a) \in R_3$  for every element  $a \in A$ .

Since  $A = \{0, 1, 2, 3\}$ ,  $R_3$  is reflexive if it contains  $(0, 0), (1, 1), (2, 2), (3, 3)$

we note that  $R_3$  does not contain  $(0, 0)$  and thus  $R_3$  is not reflexive

" $R_3$  is not reflexive"

c) The relation  $R_3$  on a Set  $A$  is symmetric if  $(b, a) \in R_3$  whenever  $(a, b) \in R_3$

we note that  $(2, 3) \in R_3$  while  $(3, 2) \notin R_3$  and there are no other ordered pairs in  $R_3$ , thus  $R_3$  is not symmetric.

" $R_3$  is not symmetric"

d) The relation  $R_3$  on a Set  $A$  is transitive if  $(a, b) \in R_3$  and  $(b, c) \in R_3$  implies  $(a, c) \in R_3$

we know that  $(2, 3) \in R_3$  and  $(3, 2) \in R_3$  while  $(2, 2) \notin R_3$  and thus  $R_3$  is not transitive

" $R_3$  is not transitive"

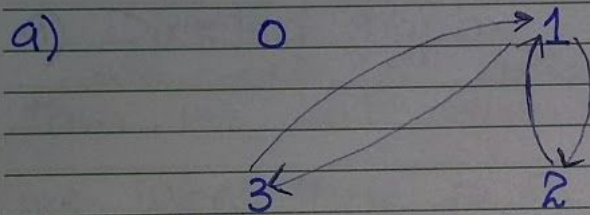
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$$R_4 = \{(1,2)(2,1)(1,3)(3,1)\}$$

$$A = \{0, 1, 2, 3\}$$

$$R_4 = \{(1,2)(2,1)(1,3)(3,1)\}$$



b) No,  $R_4$  is not reflexive because there is no loop at 1.

" $R_4$  is not reflexive"

c) ~~No~~ Yes,  $R_4$  is symmetric because between two numbers where there exist an arrow, the arrows go both ways.

" $R_4$  is symmetric"

d) No,  $R_4$  is not transitive though there are arrows between 2 and 1 and also arrows between 1 and 3 but there are no arrows between 2 and 3.

" $R_4$  is not transitive"

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$$R_s = \{(0,0)(0,1)(0,2)(1,2)\}$$

$$A = \{0, 1, 2, 3\}$$

$$R_s = \{(0,0)(0,1)(0,2)(1,2)\}$$

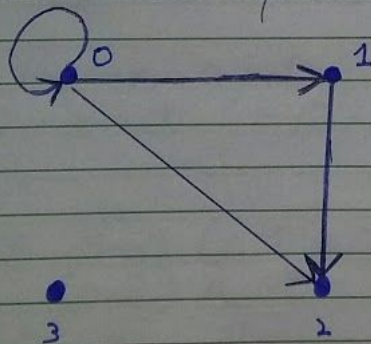
a) Directed graph:-

we note  $A$  contains 4 elements and thus we will draw 4 points.

we label these points 0, 1, 2, 3 (which are the elements of  $A$ .)

For every element  $(x,y) \in R$  with  $x \neq y$ , we draw an arrow from  $x$  to  $y$ .

For every element  $(x,x) \in R$ , we draw a loop at the point  $x$ .



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b) The Relation  $R_S$  is reflexive if  $(a,a) \in R_S$  for every element  $a \in A$ .

Since  $A = \{0, 1, 2, 3\}$ ,  $R_S$  is reflexive if it contains  $(0,0) \cup (1,1) \cup (2,2) \cup (3,3)$

we note  $R_S$  does not contain  $(1,1)$  and thus  $R_S$  is not reflexive.

" $R_S$  is not reflexive"

c) The  $R_S$  on a Set  $A$  is Symmetric if  $(b,a) \in R_S$  whenever  $(a,b) \in R_S$

we note that  $(0,1) \in R_S$  while  $(1,0) \notin R_S$  and thus  $R_S$  is not symmetric

" $R_S$  is not symmetric"

d) The  $R_S$  on a set  $A$  is Transitive if  $(a,b) \in R_S$  and  $(b,c) \in R_S$  implies  $(a,c) \in R_S$

we note that  $(0,1) \in R_S$  and  $(1,2) \in R_S$  while  $(0,2) \in R_S$  and thus  $R_S$  is transitive as there are no other pairs for which transitivity can be checked

" $R_S$  is transitive"

SMNotes

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$$R_s = \{(0,1), (0,2)\}$$

a) Directed graph:-

$$A = \{0, 1, 2, 3\}$$

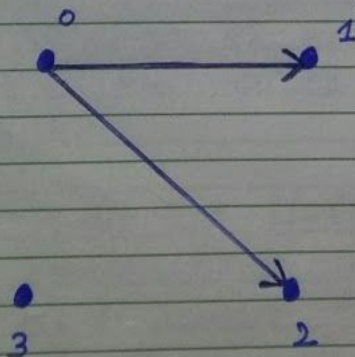
$$R_s = \{(0,1), (0,2)\}$$

we note that  $A$  contains 4 elements and thus we will draw 4 points

we label these points 0, 1, 2, 3 (which are the elements of  $A$ )

For every element  $(x,y) \in R_s$  with  $x \neq y$ , we draw an arrow from  $x$  to  $y$ .

For every element  $(x,x) \in R_s$  we draw a loop at the point  $x$ .



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2)

b) The relation  $R_6$  is reflexive if  $(a,a) \in R_6$  for every element  $a \in A$ .

Since  $A = \{0, 1, 2, 3\}$ ,  $R_6$  is reflexive if it contains  $(0,0), (1,1), (2,2), (3,3)$

we note that  $R_6$  does not contain  $(0,0)$  and thus  $R_6$  is not reflexive.

" $R_6$  is not reflexive"

c) The relation  $R_6$  on a set  $A$  is symmetric if  $(b,a) \in R_6$  whenever  $(a,b) \in R_6$

we note that  $(0,1) \in R_6$  while  $(1,0) \notin R_6$  and thus  $R_6$  is not symmetric.

" $R_6$  is not symmetric"

d) The relation  $R_6$  on a set  $A$  is transitive if  $(a,b) \in R_6$  and  $(b,c) \in R_6$  implies  $(a,c) \in R_6$

we note that the if-statement  $(a,b) \in R_6$  and  $(b,c) \in R_6$  is never true (as there exists no such pairs in  $R_6$ ). when the if-statement is false and thus  $R_6$  is transitive.

" $R_6$  is ~~not~~ transitive"

SMNotes



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$$R_7 = \{(0,3), (2,3)\}$$

$$A = \{0, 1, 2, 3\}$$

$$R_7 = \{(0,3), (2,3)\}$$

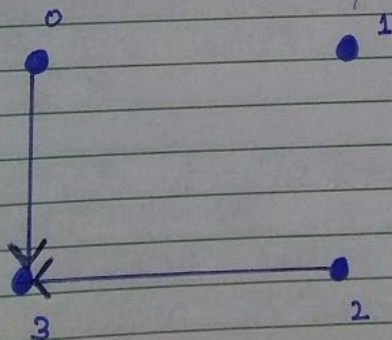
a) Directed graph:-

we note that  $A$  contains 4 elements and thus we will draw 4 points

we label these points 0, 1, 2, 3 (which are the elements of  $A$ ).

For every element  $(x,y) \in R$   $x \neq y$ , we draw an arrow from  $x$  to  $y$ .

For every element  $(x,x) \in R$ , we draw a loop at the point  $x$ .



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b) The relation  $R_7$  is reflexive if  $(a,a) \in R_7$  for every element  $a \in A$ .

Since  $A = \{0, 1, 2, 3\}$ ,  $R_7$  is reflexive if it contains  $(0,0), (1,1), (2,2), (3,3)$

we note that  $R_7$  does not contain  $(0,0)$  and thus  $R_7$  is not reflexive.

" $R_7$  is not reflexive"

c) The  $R_7$  on a set  $A$  symmetric if  $(b,a) \in R_7$  whenever  $(a,b) \in R_7$

we note that  $(0,3) \in R_7$  while  $(3,0) \notin R_7$  and thus  $R_7$  is not symmetric

d) " $R_7$  is not symmetric"

d) The  $R_7$  on a set  $A$  is transitive if  $(a,b) \in R_7$  and  $(b,c) \in R_7$  implies  $(a,c) \in R_7$

We note that if-statement  $(a,b) \in R_7$  and  $(b,c) \in R_7$  is never true as there exist no such pairs in  $R_7$ . However an if-then statement is always true

~~when the~~ " $R_7$  is transitive"

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$$R_8 = \{(0,0) \text{ \& } (1,1)\}$$

$$A = \{0,1,2,3\}$$

$$R_8 = \{(0,0), (1,1)\}$$

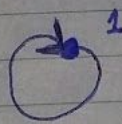
a) Directed graph:-

we note that  $A$  contains 4 elements and thus we will draw 4 points.

we label these points 0, 1, 2, 3 (which are the elements of  $A$ ).

For every element  $(x,y) \in R$  with  $x \neq y$ , we draw an arrow from  $x$  to  $y$ .

For every element  $(x,x) \in R$ , we draw a loop at the point  $x$ .



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b) The relation  $R_8$  is reflexive if  $(a,a) \in R_8$  for every element  $a \in A$ .

Since  $A = \{0, 1, 2, 3\}$   $R_8$  is reflexive if it contains  $(0,0) \cup (1,1) \cup (2,2) \cup (3,3)$

we note that  $R_8$  does not contain  $(2,2)$  and thus  $R_8$  is not reflexive.

" $R_8$  is not reflexive"

c) The relation  $R_8$  on a set  $A$  is symmetric if  $(b,c) \in R_8$  whenever  $(a,b) \in R_8$

we note that  $(0,0) \in R_8$  while  $(0,0) \in R_8$  and  $(1,1) \in R_8$  while  $(1,1) \in R_8$ . There are no other elements in  $R_8$  so the  $R_8$  is symmetric.

" $R_8$  is symmetric"

d) The  $R_8$  on a set  $A$  is ~~symmetric~~ transitive if  $(a,b) \in R_8$  and  $(b,c) \in R_8$  implies  $(a,c) \in R_8$

we note that the statement if " $(a,b) \in R_8$  and  $(b,c) \in R_8$ " is never true for two distinct pairs (as there no other such pairs in  $R_8$ ). when if statement is false so  $R_8$  is transitive.

" $R_8$  is transitive"

MASTER NOTES

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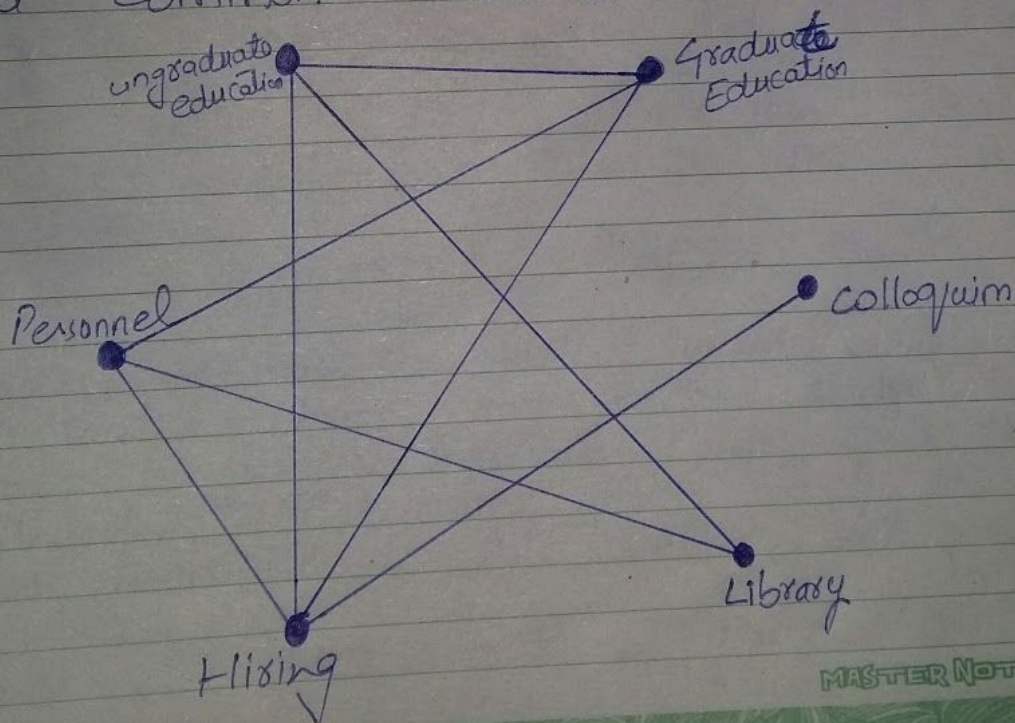
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"Question no 10"

Answer:-

We have been given 6 Committees: undergraduate Education, Graduate education, Colloquium, Library, Hiring and personnel. Let us then represent the scheduling problem by a graph with 6 vertices and each vertex represents one of the committees.

Next, we draw an edge between two vertices, if the two committees have a common member.



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Next, we will assign a color to each other vertex such that no two adjacent vertices have the same color (thus no two vertices that are connected by an edge will have the same color).

we note that the vertex "Hiring" has the most edges connecting the vertex to another vertex. Thus let us color this hiring vertex first the color Red.

"Hiring" = Red

Next, the vertex "Library" is the only vertex that is not connected to the vertex "Hiring" by an edge and thus we can color "Library" the same as "Hiring". Let us color the vertex "Library" with the color ~~purple~~

"Library" = ~~purple~~ Red

Next, we note the vertex "Personnel" is connected to vertex "Hiring" and vertex ~~Library~~ "Library" by an edge. Thus we cannot color the vertex "Personnel" with Red, let us then color this vertex with purple.

"Personnel" = purple.

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Next, we note vertex "Graduate Education" is connected to vertex "Hiring" and vertex "personnel" by an edge. Thus we cannot color the vertex "personnel" with Red nor with purple, let us then color this vertex with blue.

"Graduate Education" = blue

Next, we note that vertex "undergraduate education" is connected to vertex "Hiring" and vertex "Graduate education" by an edge. Thus we cannot color the vertex "personnel" with Red nor with blue, let us then color this vertex with purple.

"undergraduate education" = purple.

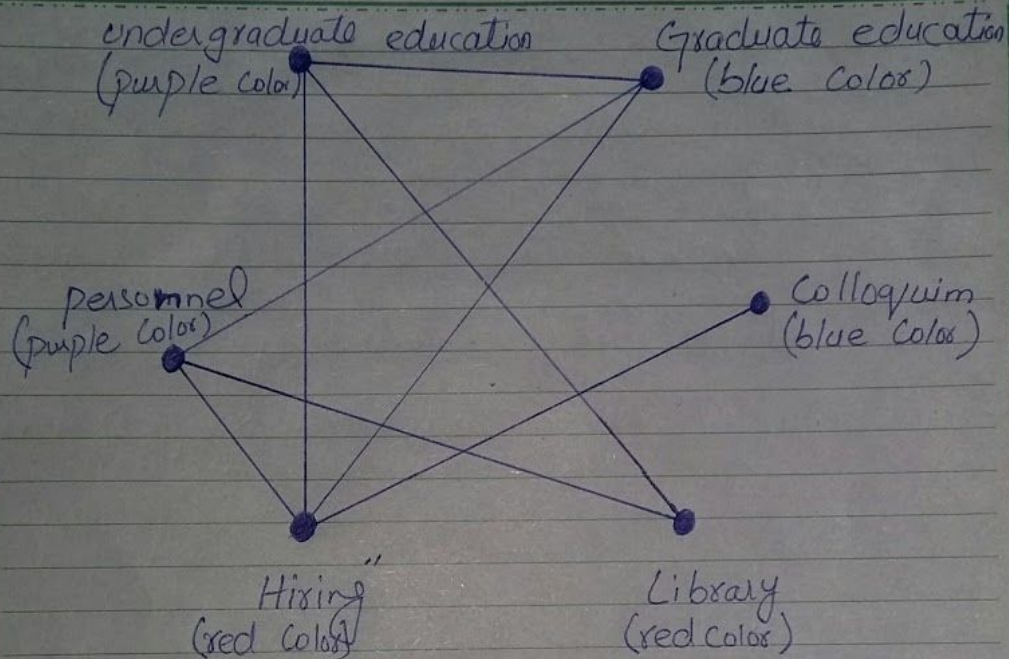
And then connect the Colloquium with the color blue

"Colloquium" = blue

we then note that we colored all vertices such that no adjacent vertices have the same color and we used exactly 3 colors.

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we can then schedule the meeting by assigning each time slot to a color.

Thus let us assign the first time slot to the red color, the second time slot to the blue color and the third time slot to the purple color.

First time slot: Hiring and Library

Second time slot: colloquium and Graduate education

Third time slot: undergraduate education and personnel.