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Q1:- Explain the following operation in a single link list.

- 1) Insert an Element
- 2) Delete an Element

Answer:-

Now that you have got an understanding of the basic concepts behind linked list and their types its time to dive into the common operations that can be performed.

Two important points to remember.

- ⇒ (head) points to the first node of the linked list
- ⇒ pointer of the last node is (Null) so if the next current node is (Null) we have reached the end of the linked list.

In all of the example, we will assume that the linked list has three nodes 1 → 2 → 3 with node structure as below.

```

struct node
{
    int data;
    struct node * next;
};
    
```

⇒ How To Traverse a linked list:-

When temp (Null) we know that we have reached the end of the linked list so we get out of the while loop.

```

struct node * temp = head;
printf ("%d\n", list element are -> {n});
while (temp != Null)
{
    printf ("%d -> ", temp->data);
    temp = temp->next;
}

```

output:

list element are

1 --> 2 --> 3 -->

⇒ How to Add Element to a linked list

⇒ Add to the beginning:

- ① Allocate memory for new node
- ② store data
- ③ change next of new to point to head
- ④ change head to point to recently created node

```

struct node * newNode;
newNode = malloc (size of (struct node));
newNode->data = 4;
newNode->next = head;
head = newNode;

```

→ Add to the End:-

- ① Allocate memory for new node
- ② Store data
- ③ Traverse to last node
- ④ Change next of last node to recently created node.

```

struct node * newnode;
newnode = malloc (sizeof (struct node));
newnode -> data = 4;
newnode -> next = Null;
struct node * temp = head;
while (temp -> next != Null) {
    temp = temp -> next;
}
temp -> next = newnode;

```

→ Add to the Middle

- ① Allocate memory and store data for new node
- ② Traverse to node just before the required position of new node
- ③ change next pointer to include new in b/w.

```

struct node * newNode;
newNode = malloc (sizeof (struct node));
newnode -> data = 4;
struct node * temp = head;
for (int i = 2; i < position; i++) {
    if (temp -> next != Null) {
        temp = temp -> next;
    }
}

```

newnode \rightarrow next = temp \rightarrow next
 temp \rightarrow next = newnode;

\Rightarrow How to Delete:-

\Rightarrow Delete from beginning

② Point head to the second node

`head = head \rightarrow next;`

\Rightarrow Delete from end:

① Traverse to second last element

② change its next pointer to null.

struct node* temp = head;

while (temp \rightarrow next \rightarrow next \neq Null) {

temp = temp \rightarrow next;

}
 temp \rightarrow next = Null

\Rightarrow Delete from middle:-

① Traverse to element before the element to be deleted

② change next pointer to exclude the node from the chain.

for (int i=2; i < position; i++) {

if (temp \rightarrow next \neq Null) {

temp = temp \rightarrow next

}

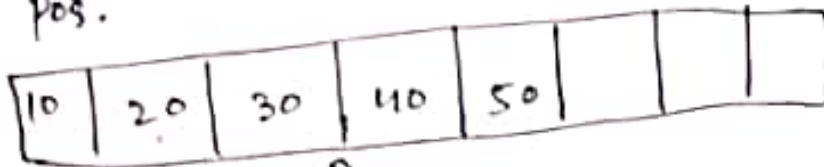
}

temp \rightarrow next = temp \rightarrow next \rightarrow next;

Q2: write a program to insert a new element in the given unsorted array at K^{th} position

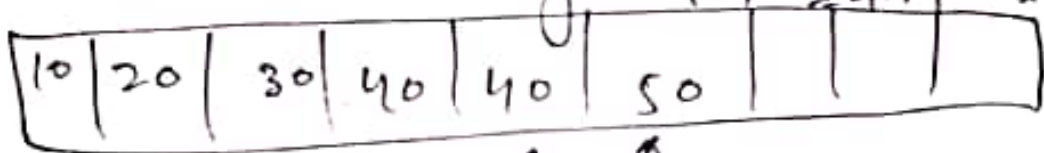
Answer:

- ① Input size and elements in array. Store it in some variable say size and arr.
- ② Input new element and position to insert in array. Store it in some variable say num and pos.

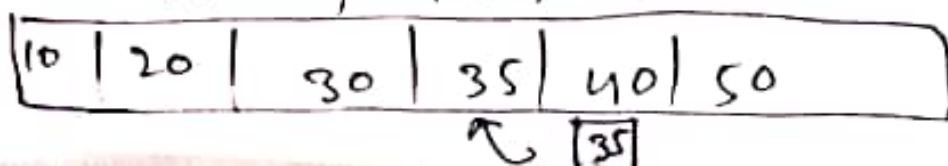


- ③ To insert new element in array shift element from the given position to one position right hence run a loop in descending order from size to pos to insert. The loop structure should look like for ($i = \text{size}; i > \text{pos}; i--$)

inside the loop copy previous element to current element by $\text{arr}[i] = \text{arr}[i-1]$;



- ④ finally after performing shift operation. copy the new element at its specified position i.e. $\text{arr}[\text{pos}-1] = \text{num}$;



⇒ program to insert element in array:

```

/**
 * C program
 */
#include <stdio.h>
#define Max_size 100

int main()
{
    int arr[Max_size];
    int i, size, num, pos;

    /* input size of the array */
    printf("Enter size of the array: ");
    scanf("%d", &size);

    /* input element in array */
    printf("Enter element in array: ");
    for (i=0; i << size; i++)
    {
        scanf("%d", &arr[i]);
    }

    /* input new element and position to insert */
    printf("Enter element to insert: ");
    scanf("%d", &num);
    printf("Enter the element position: ");
    scanf("%d", &pos);

    /* if position of element is not valid */
    if (pos > size+1 || pos <= 0)
    {
        printf("Invalid position, please enter position b/w 1 to %d", size);
    }
    else

```

/* make room for new array element
by shifting to right */

```
for (i = size; i >= pos; i--)  
{  
  arr[i] = arr[i-1];  
}
```

/* insert new element at given position and
increment size */ arr[pos-1] = num;
size++;

/* print array after insert operation */

```
print ("Array element after insertion:");  
{  
  for (i=0; i < size; i++)  
    printf ("%d\t", arr[i]);  
}  
return 0;
```

}



Q3:-

Answer:

Quick Sort:-

Quick Sort is a fast sorting algorithm used to sort a list of elements. Quick sort algorithm is invented C.A.R Hoare.

The quick sort algorithm attempts to separate the list of element into two parts and then sort each part recursively. that means

it use divid and conquer strategy. In quick sort the position partition of the list is performed based on the element called pivot. Here pivot element is one of the elements in the list.

The list is divided into two partition such that "all element to the left are smaller than the pivot and all elements to the right of pivot are greater than or equal to the pivot."

Example

consider the following unsorted list elements.

list

5	3	8	1	4	6	2	7
---	---	---	---	---	---	---	---

Define pivot, left and right. Set pivot = 0, left = 2 and right = 7. Here 7 indicates size - 1.

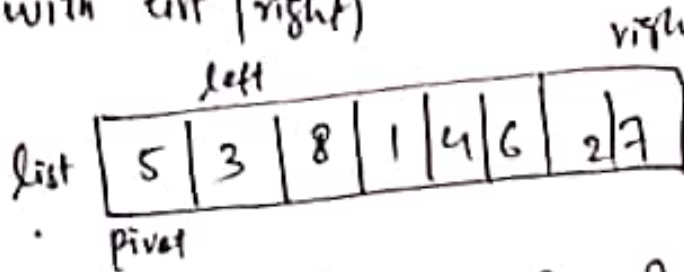
list

5	3	8	1	4	6	2	7
---	---	---	---	---	---	---	---

pivot

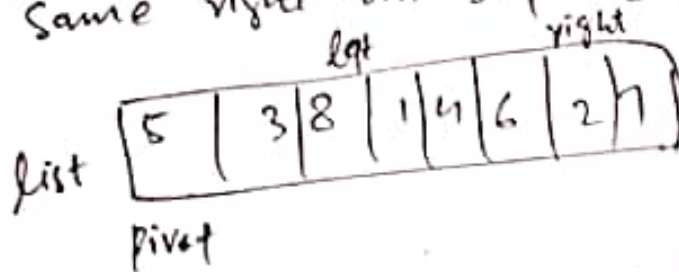
compare list [left] with list [pivot]. If list left is greater than list [pivot] then stop. left otherwise move left to the next
compare list [right] with list [pivot]. If list [right] is smaller than list [pivot] then stop right otherwise move right to the previous.

Repeat the same until $left \geq right$
 If both left & right are stopped but $left < right$
 then swap $list(left)$ with $list(right)$ and continue
 the process if $left \geq right$ then swap $list(pivot)$
 with $list(right)$

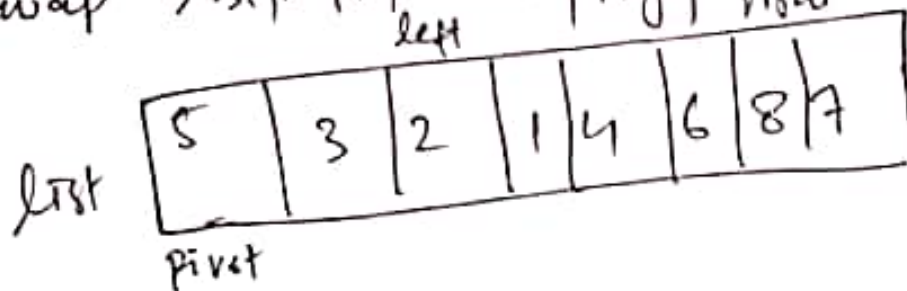


compare $list(left) < list(pivot)$ as it is true
 increment left by one and repeat the same
 left will stop at 8.

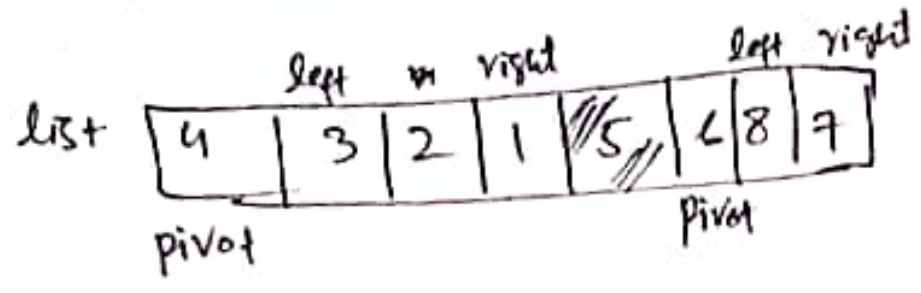
compare $list(right) > list(pivot)$ as it is true
 decrement right by one and repeat the
 same right will stop at 2



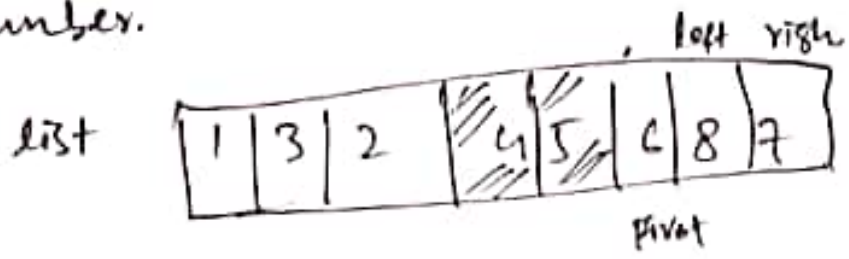
Here left & right both are stopped and left is
 not greater than right. so we need to
 swap $list(left)$ and $list(right)$



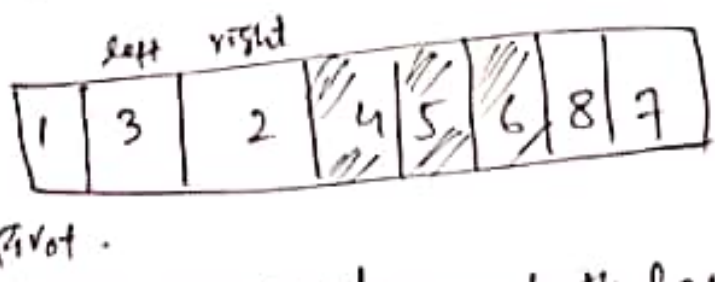
and then



In the left sublist as there are no smaller number than the pivot left will keep on moving to the next and stop at last number.



and then



Repeat the same recursively on both left and right sublists until all the numbers are sorted. The final sorted list will be as follows.



⇒ Insertion Sort Algorithm:

a : array of items

n : size of list

Start

1. declare variable - i, key, j
 2. loop: $i = 1$ to $n-1$ // outer loop
 - 2.2 $\text{key} = a[i]$
 - 2.2 $j = i - 1$
 - 2.3 loop: ($j >= 0$ & $a[j] > \text{key}$) // inner loop
 - 2.3.1 $a[j+1] = a[j]$
 - 2.3.2 $j = j - 1$
 - end loop // inner loop
 - 2.4 $a[j+1] = \text{key}$
 - end loop // outer loop
- Stop.

Q4:-



Answers:-

⇒ part(A) Adjacency matrix:-

	A	B	C	D	E
A	0	1	0	1	0
B	1	0	1	0	0
C	0	1	0	1	1
D	1	0	1	0	1
E	0	0	1	1	0

1 = true, 0 = false

Need $O(N^2)$ space to store the graph
with sparse graphs (graphs that do not have
too many edges), there may be "too many"
size zeros in the matrix. Dense graphs
have many more edges.

→ Part (B)

