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Department

BS(CS)

Semester

2th

Subject

Discreat sturture

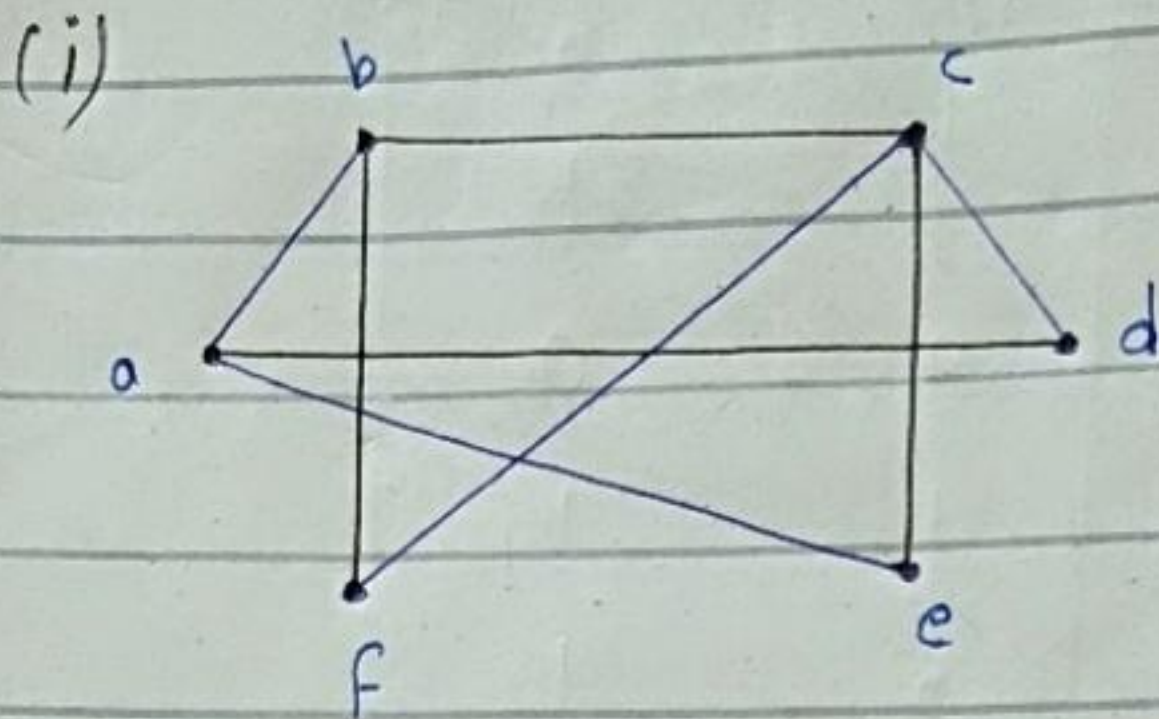
Sumbi

Submitted to : Saif Ullah Jan

Date : 20/6/2020

Assignment : 3

# Question 1



This is NOT bipartite.

**Bipartite Graph**  $\Rightarrow$

A bipartite graph is a simple graph whose vertices can be partitioned into set  $V_1$  and  $V_2$  such that there are no edges among the vertices of  $V_1$  and no edges among the vertices of  $V_2$ , while there can be edges between a vertex of  $V_1$  and vertex of  $V_2$ .

A simple graph is bipartite if and only if it is possible to assign red or blue to every vertex such that no two connected vertices have the same color.

Let us assign red or blue to each vertex.

If two vertices are connected, then they should not have the color.

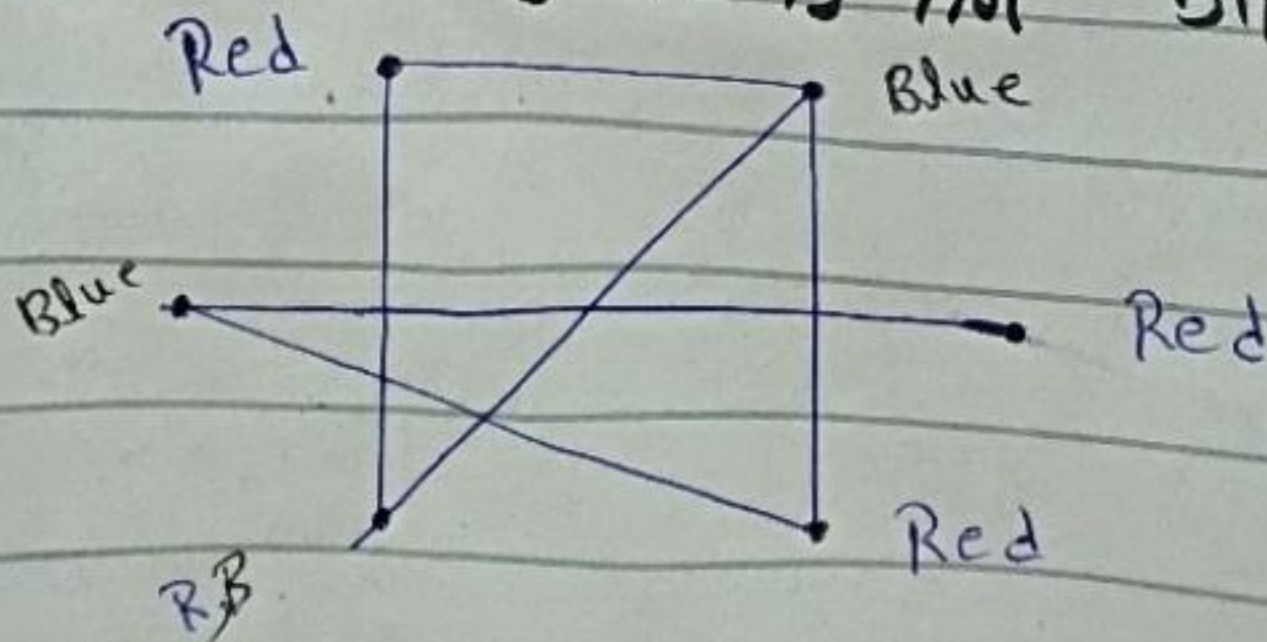
We start by assigning "Blue" to a.

- Since b is connected to the blue a, we assign "Red" to b.

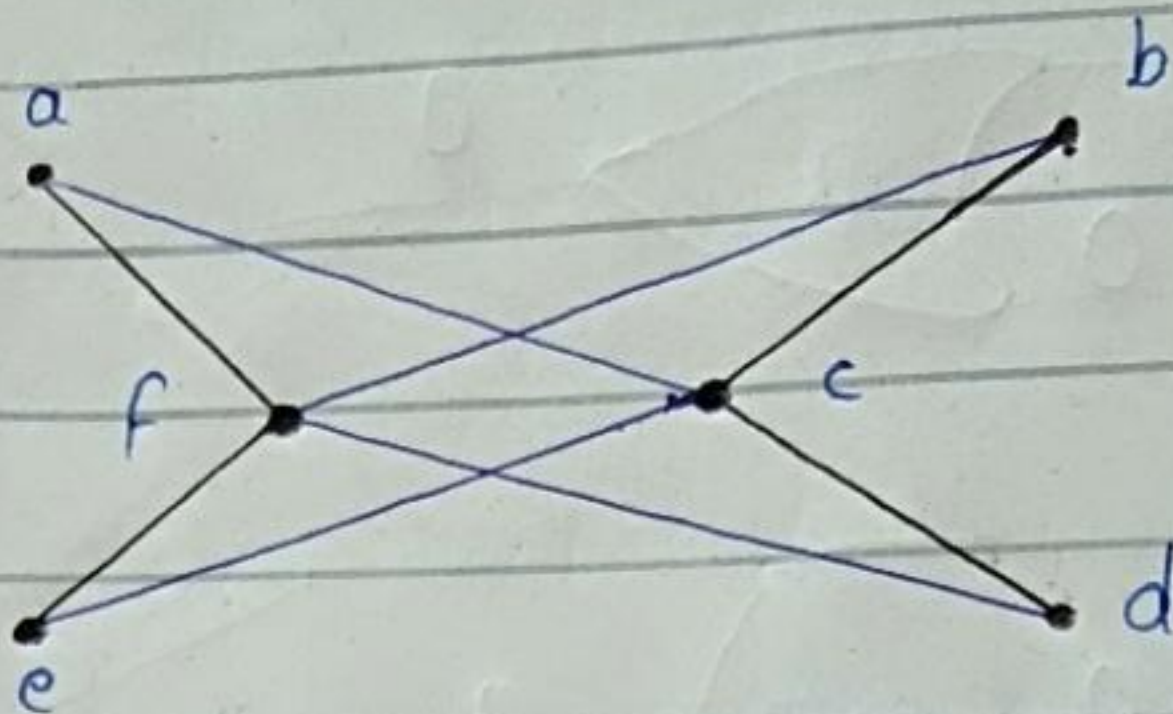
- Since c is connected to the red b, we assign "Blue" to c.

We then note that f is connected to the red b and the blue c, which means that we cannot assign a color to f such that it differ from the color of the connect vertices. (As there are only two color: Red/blue)

Thus it is not possible to assign red or blue to each vertex such that connect vertices do not have the same color the graph is not bipartite.



(ii)

Solution  $\Rightarrow$ 

This is bipartite.

A simple graph is bipartite if and only if it is possible to assign red or blue to every vertex such that no two connected vertices have the same color.

Let us assign red or blue to each vertex.

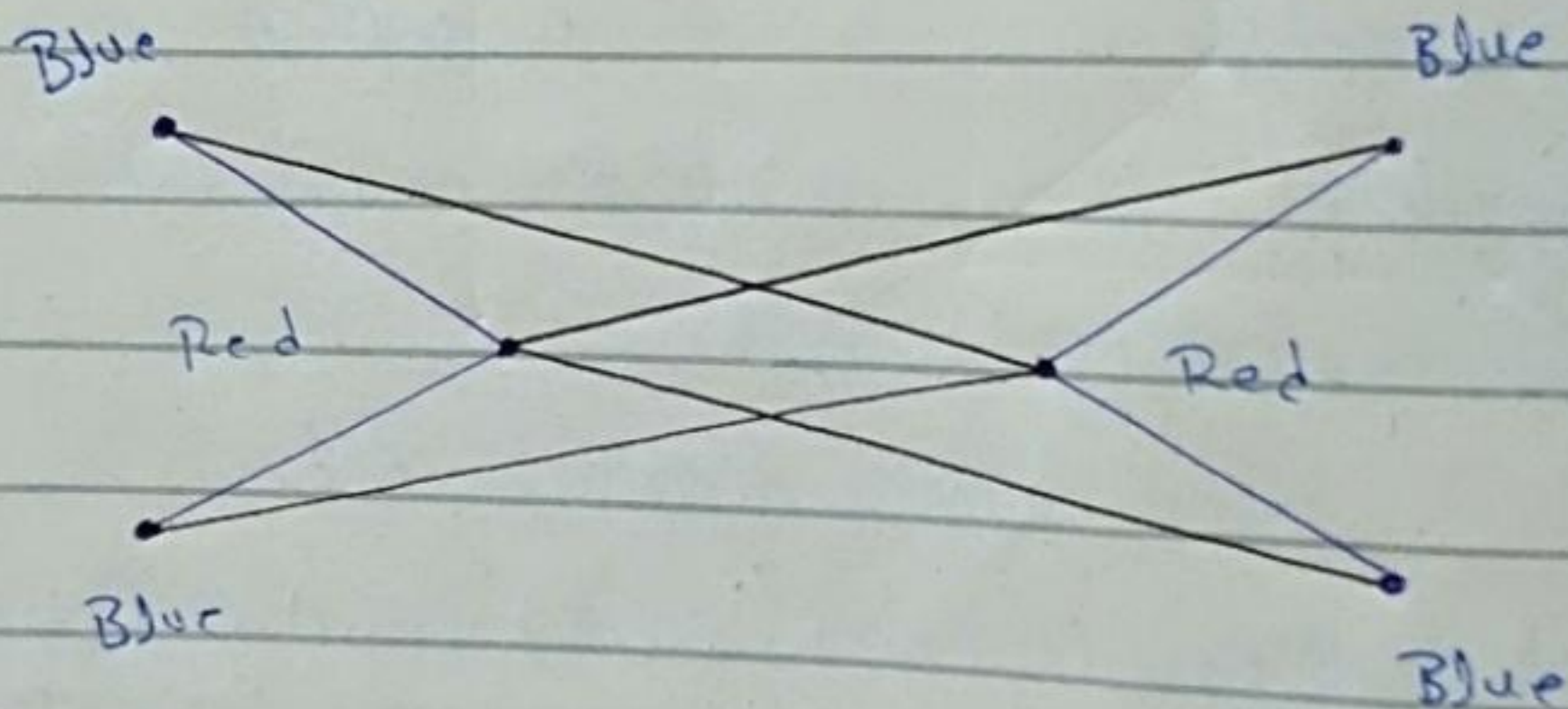
If two vertices are connected, then should not have the same color. We then note that it is possible to assign red or blue to each vertex such that connected vertices do not have the same color and thus the

graph is bipartite.

Moreover, the Partitioning of the vertices and the set with the red vertices

$$V_1 = \{a, b, c, d\}$$

$$V_2 = \{c, f\}$$

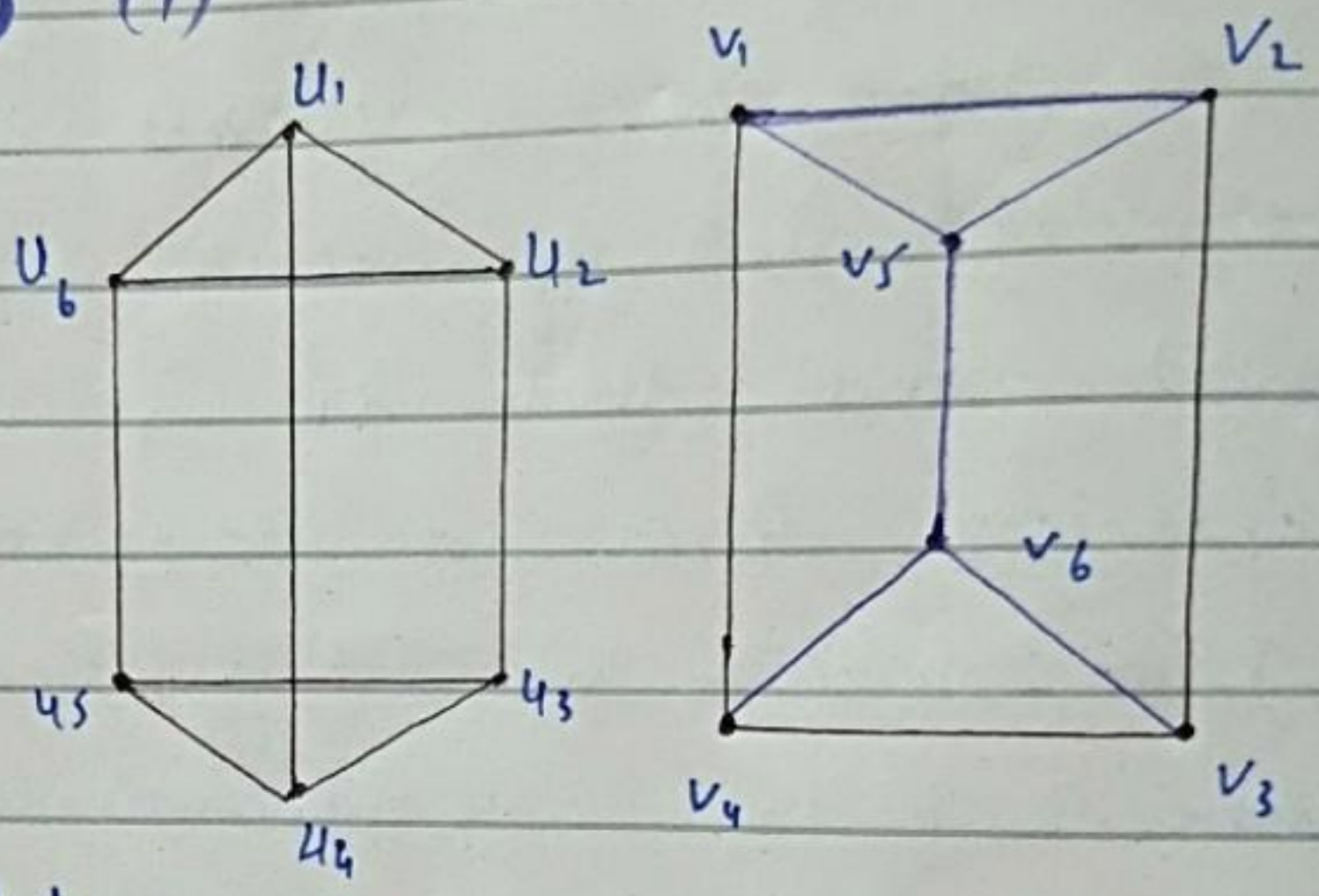


Graph is Bipartite.

# Question 2

Graph is Isomorphic

(a) (i)



Solution

Let us first determine set of vertices of edge.

$$V_1 = \{ u_1, u_2, u_3, u_4, u_5, u_6 \}$$

$$E_1 = \{ (u_1, u_2), (u_1, u_6), (u_2, u_6), (u_2, u_3), (u_6, u_5), (u_3, u_5), (u_3, u_4), (u_5, u_4), (u_1, u_4) \}$$

Let us the first determine set of vertices and set of edge right graph.

$$V_2 = \{ v_1, v_2, v_3, v_4, v_5, v_6 \}$$

$$E_2 = \{ (v_5, v_2), (v_5, v_4), (v_6, v_1), (v_2, v_3), (v_2, v_1), (v_3, v_4), (v_3, v_6), (v_4, v_6), (v_4, v_5) \}$$

By computing the two set of edge we can define the following

one to one and onto function  
f from  $V_1$  to  $V_2$ .

note: you could also use the degree of the vertices because and their image need to have same degree.

$$f(u_1) = v_5 \quad f(u_2) = v_2$$

$$f(u_3) = v_3 \quad f(u_4) = v_6$$

$$f(u_5) = v_4 \quad f(u_6) = v_1$$

f is then a function that make the two graph isomorphic

$u_1$  and  $u_2$  are adjacent  $f(u_1) = v_5$  and  $f(u_2) = v_2$  are adjacent  
 $u_1$  and  $u_4$  are adjacent, while  $f(u_1) = v_3$  and  $f(u_4) = v_6$  are adjacent.  
 $u_1$  and  $u_6$  are adjacent while  $f(u_1) = v_5$  and  $f(u_6) = v_1$

~~This is isomorphic.~~

are adjacent.

$u_2$  and  $u_3$  are adjacent, while  $f(u_2) = v_2$  and  $f(u_3) = v_3$  are adjacent.

$u_3$  and  $u_4$  are adjacent, while  $f(u_2) = v_2$  and  $f(u_6) = v_1$  are adjacent.

$u_3$  and  $u_5$  are adjacent, while,  $v_6$  and  $f(u_5) = v_4$  are adjacent.

$u_5$  and  $u_6$  are adjacent, while  $f(u_5) = v_4$  and  $f(u_6) = v_1$  are adjacent.

isomorphic.





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$f$  is then a function  
that make two graph  
isomorphic.

This is also isomorphic.

# Question 3

(a)

$$= \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

Solution:

Let us add all element in the two matrices, which will represent the number of connection of a vertex to another vertex which is double the number edge.

A contains 8 ones and thus the simple graph corresponding to A contain 8 connection.

B contains 10 ones and thus the simple graph corresponding to A contain 10 connections.

Since the number of connection of the two graph are not the same, the number of edges in the graph are not the same and then the Graph are not isomorphic.

b

$$= \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix}$$

Solution:

Let us add all element in the two matrices, which will represent the number of connection of a vertex to another vertex.

A contain 8 ones and thus the simple graph corresponding to A contain 8 connection.

B contain one and thus the simple graph corresponding to A contain 6 connection.

Since the number of connection of the two graph are not the same, the number of edges in the graph are not the same and then the Graph are not isomorphic.

## Question 4

(a)

Euler circuit  $\Rightarrow$ 

An Euler circuit is a simple circuit that contains every edge of the graph.

A path in a directed graph  $G$  is a sequence of edges in  $G$ .

A simple path is a path that does not contain the same edge more than once. A circuit is a path that begins and ends in the same vertex.

The degree of a vertex is the number of edges that connect to the vertex.

Let us first determine the degree of every vertex in the given graph.

$$\text{deg}(a) = 2$$

$$\text{deg}(b) = 4$$

$$\text{deg}(c) = 2$$

$$\text{deg}(d) = 4$$

$$\text{deg}(e) = 4$$

$$\text{deg}(f) = 4$$

$$\text{deg}(g) = 2$$

$$\text{deg}(h) = 4$$

$$\text{deg}(i) = 2$$

A graph has an Euler circuit if and only if each of the vertices has an even degree. Since

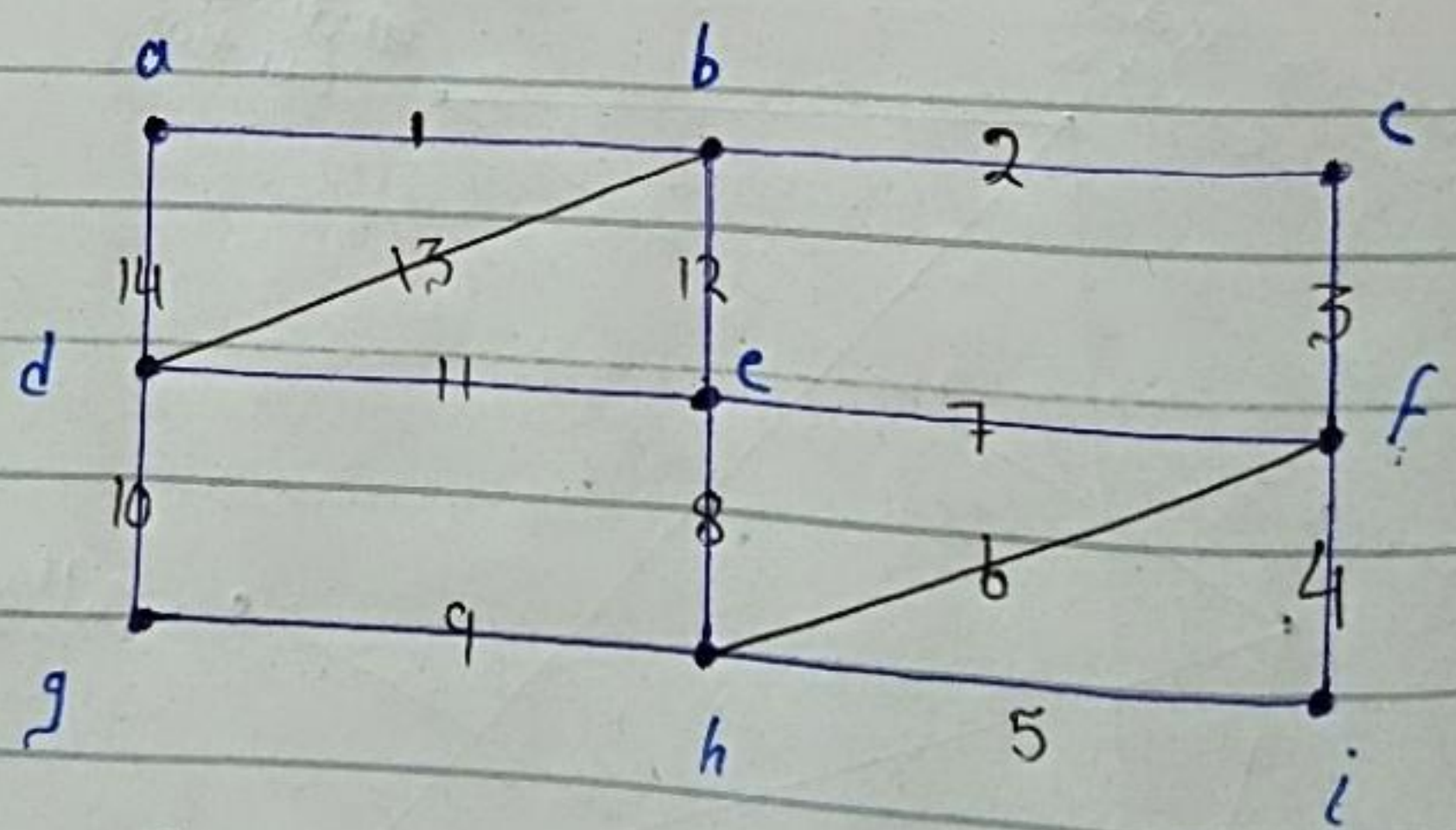
degree are even, there exist  
an Euler circuit.

$\text{deg}(g) = 2$        $\text{deg}(h) = 4$

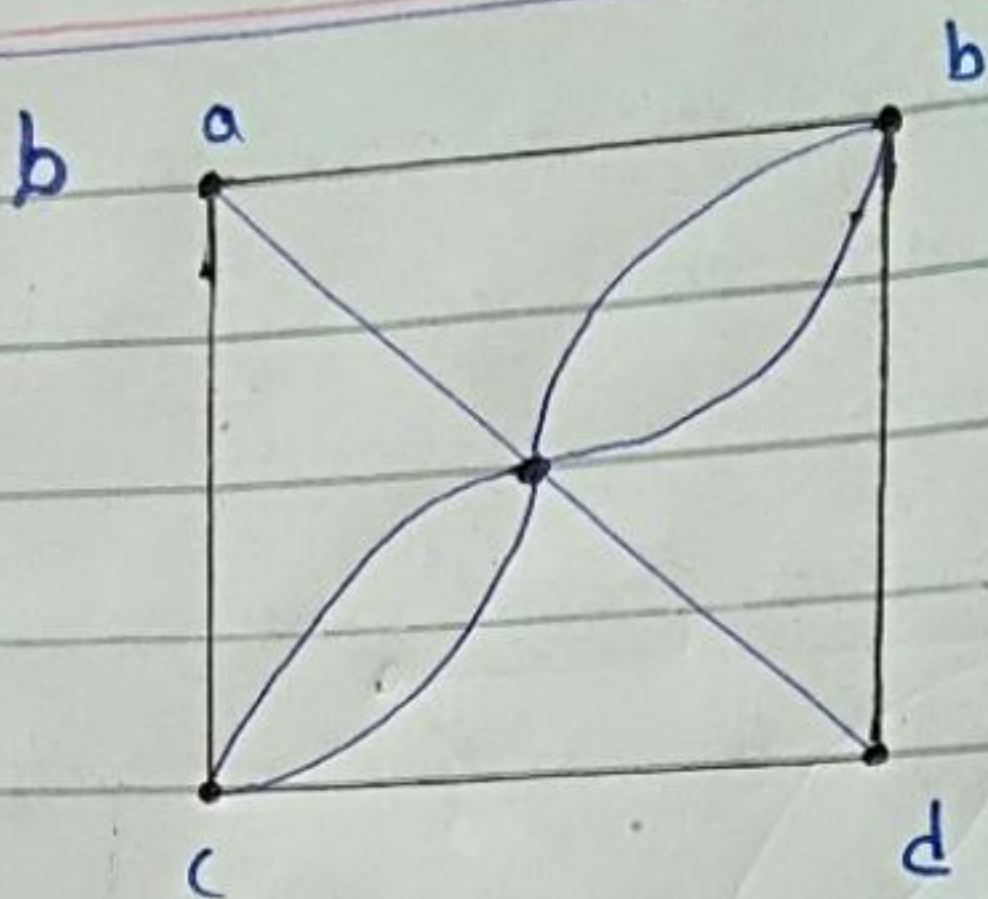
$\text{deg}(i) = 2$

A graph has Euler circuit if  
and only if each of the vertices  
has an even degree. Since all  
degree are even, there exist Euler circuit.

A Possible Euler circuit is  
a, b, c, f, i, h, f, c, h, g, d, c, b, d, a



This is Euler circuit exist.



Solution:

Let us first determine the degree of every vertex in the given graph.

$$\text{deg}(a) = 3 \quad \text{deg}(b) = 4$$

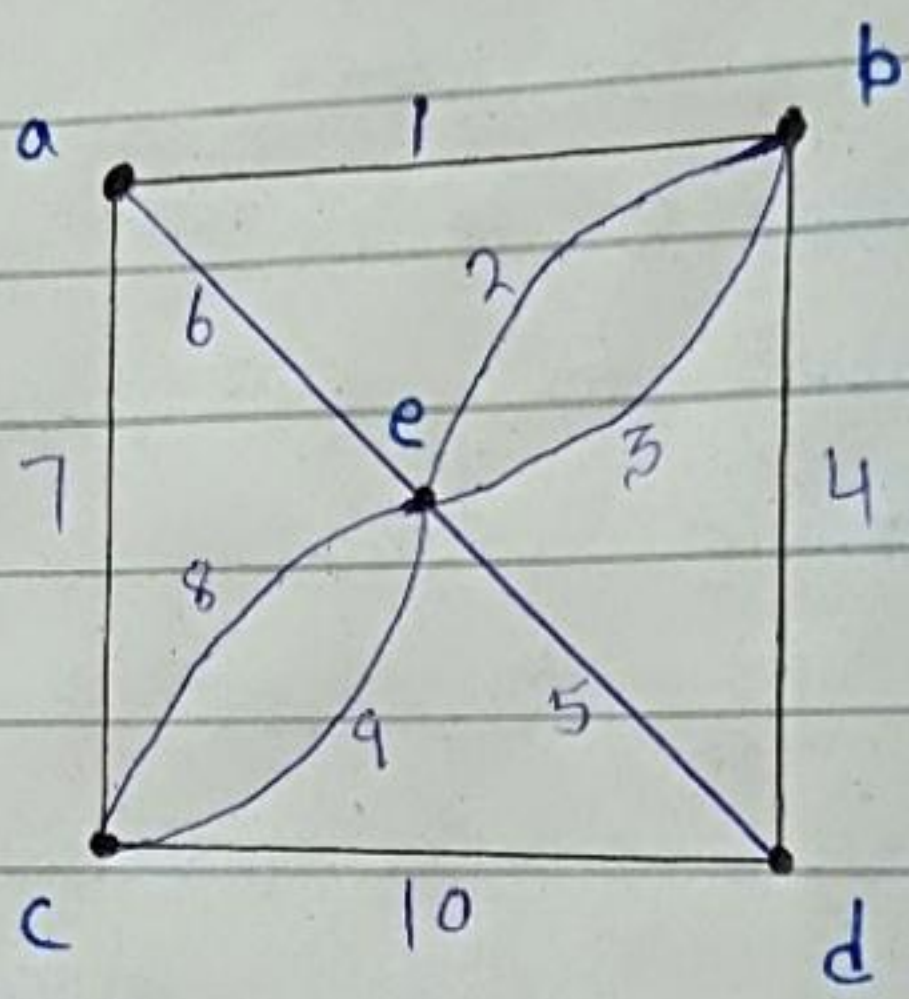
$$\text{deg}(c) = 4 \quad \text{deg}(d) = 3$$

$$\text{deg}(e) = 6$$

A graph has an Euler circuit if and only if each of the vertices has an even degree. Since some degrees are odd, there is no Euler circuit.

A graph has an Euler path if and only if there are exactly two vertices with an odd degree. We note that vertices a and d have an odd degree. Thus an Euler path exists.

A Possible Euler Path is  
a, b, c, ~~d~~, b, d, e, a, c, e, e, d



No Euler circuit exists  
Euler Path exists.



## (a) Question 5. Hamilton circuit

A Hamilton circuit is a simple circuit that passes through every vertex exactly once.

Dirac's theorem A simple graph  $G$  with  $n$  vertices ( $n \geq 3$ ) has Hamilton circuit, if the degree of every vertex in  $G$  is at least  $n/2$ . A path in a directed graph  $G$  is a sequence of edge in  $G$ . A simple path is a path that does not contain the same edge more than once.

A circuit is a path that begins and ends in the same vertices.

The degree of a vertex is the number of edge that connect to the vertex. A cut edge in an edge is an edge that when removed from a graph, the resulting graph has more connected component than the original graph.

let us first determine the degree of every vertex in given

$$\deg(a) = 3$$

$$\deg(b) = 3$$

$$\deg(c) = 3$$

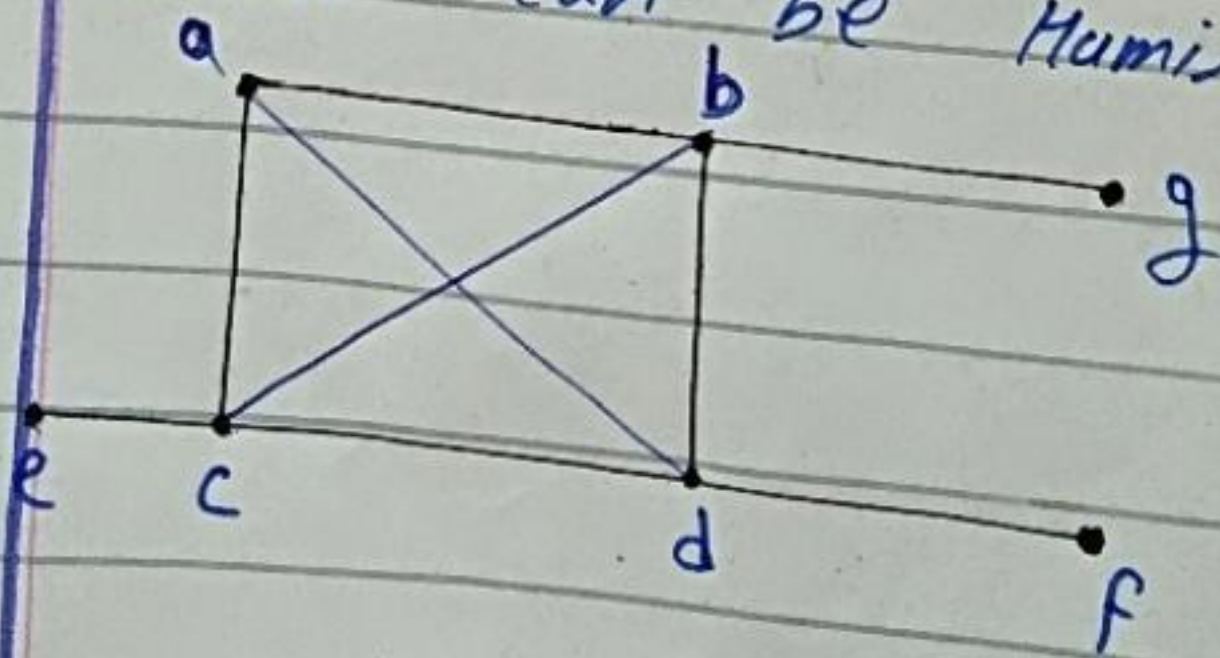
$$\deg(d) = 2$$

$$\deg(e) = 4$$

$$\deg(f) = 1$$

we then note that Dirac's theorem is not satisfied less than  $n/2 = 6/2 = 3$ , but this does not necessarily mean that no Hamilton circuit exist.

However we do note that there is only one edge  $(c, f)$  connecting to  $f$  and than any circuit that contain  $f$  need to pass through  $c$  twice, which mean than no circuit can be Hamilton circuit



Hamilton circuit does not exist because  $f$  has only 1 edge connecting it

(b)  $n = 5$

Let us first determine the degree of every vertex.

$\text{deg}(a) = 2$        $\text{deg}(b) = 4$

$\text{deg}(c) = 2$        $\text{deg}(d) = 3$

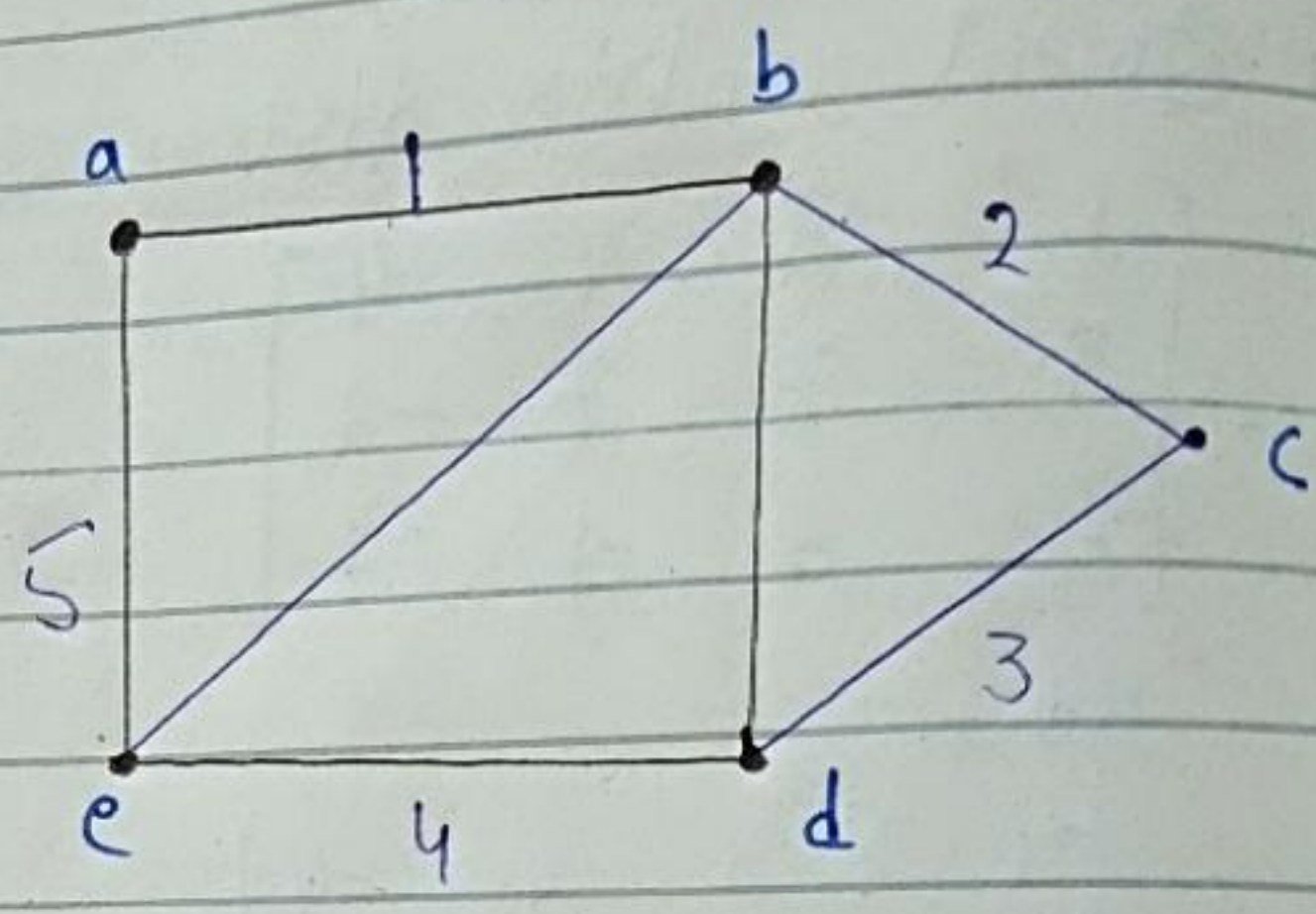
$\text{deg}(e) = 3$

We then note that Dirac's theorem is not satisfied (less than  $n/2 = 5/2 = 2.5$ ) but this does not necessarily mean that no Hamilton circuit exist

However, we do note that given graph contain the cycle  $C_5$  and the cycle  $C_5$  within the given graph form a Hamilton circuit (as the circuit will pass through all vertices exactly once).

A possible Hamilton circuit is thus the

Path  $C_5 : a, b, c, d, e, a$



This is Hamilton Circuit exist.

EMD