Department of Electrical Engineering Assignment

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Date:	07/05/2020

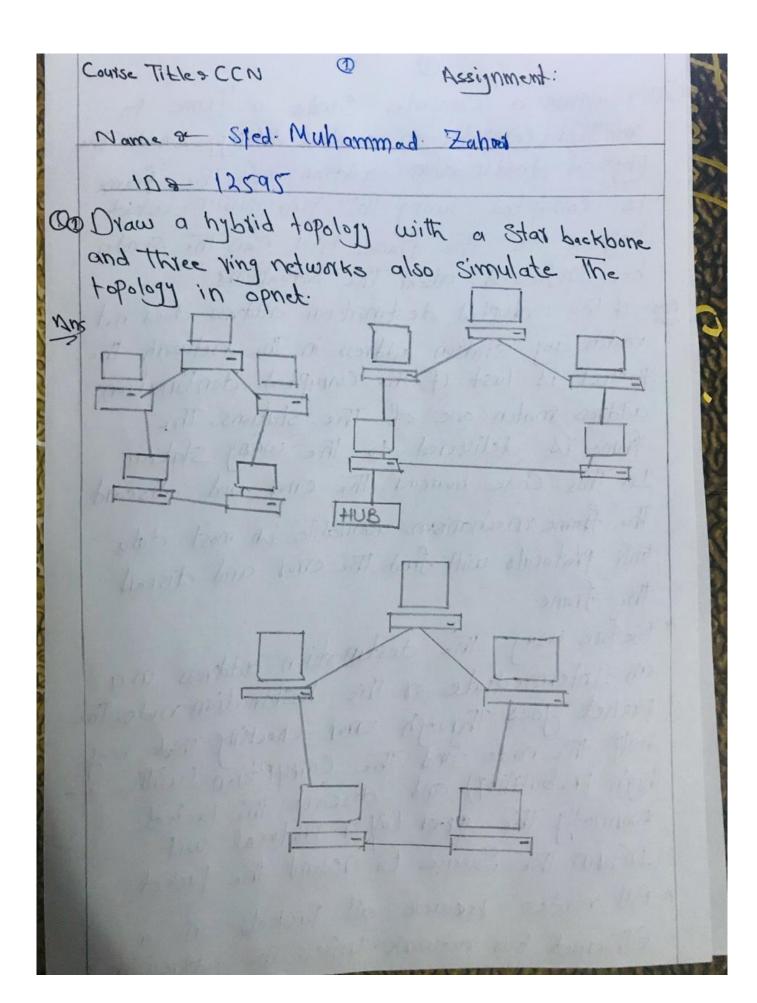
Course Details

Course Title:	Computer Communication Network	Module:	
Instructor:		Total	20
		Marke:	

Student Details

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Q1.	(a)	Draw a hybrid topology with a star backbone and three ring networks also simulate the topology	Marks 4
		in Opnet.	CLO 1
Q2.	(a)	Suppose a computer sends a frame to another computer on a bus topology LAN. The physical	Marks 4
		destination address of the frame is corrupted during the transmission. What happens to the frame?	CLO 1
		How can the sender be informed about the situation?	
Q3.	(a)	Suppose a computer sends a packet at the transport layer to another computer somewhere in the	Marks 4
		Internet. There is no process with the destination port address running at the destination	CLO 1
		computer. What will happen?	
Q4.	(a)	Match the following to one or more layers of the OSI model:	Marks 4
		a. Reliable process-to-process message delivery	CLO 1
		b. Route selection	
		c. Defines frames	
		d. Provides user services such as e-mail and file transfer	
Q5. (a)	(a)	Draw the graph of the NRZ-L, NRZ-I and Manchester scheme using each of the following data	Marks 4
		streams, assuming that the last signal level has been positive. From the graphs, guess the	CLO 2
		bandwidth for this scheme using the average number of changes in the signal level.	
		a. 00000000	
		b. 11111111	
		c. 01010101	
		d. 00110011	



Que Suppose a Computer Sends a frame to another computer on a bus toloroff LAN. The Physical destination address of the frame is corrupted during the transmission. What happens to the frame? How can the sender be Informed about the Situation?

Ans:-If the corrupted destination address does not match any station address in the network the Packet is lost. If the corrupted dentination address match one of the stations. The frame is delivered to the wrong station In this case however. The error and discard the frame mechanism, available is most data link protocals will find the error and discard the frame.

Before using The destination address in a an Intermediate or the destination node The Packet goes Through ever checking that may help the rade find the consuption (with a high probability) and discard the Packet normally the upper layer protocal will Inform the Source to resent the Packet.

+ All modes reveieve all Packets on a Ethernet bus network. Unless The address in

The Packet was not acknowledged.

(3)

Anso Transpor later is responsible for process (Source) to-process (Jestination) delivery of entire message, whereas notwork later oversees host (Source) to-host (Jestination) delivery of Endividual Packets across multiple likes The processes at each machine that communication at a given later physical later has a direct link blue I deviced while other laters have to pass the Information down to the later on the Sounder Levice by adding extra bit at each later and the received device unevrafs the message at each later moving upwards the message at each later moving upwards the finally reached the corresponding communicating later.

blu devices. At the higher lapers however communication must move down through the layers on Sending device over to receiving device, and then back up through the layers fach layer in the Sending device adds its own Information to the message it reverences from the layer Just above it and passes the whole package to the layer Just below it At layer I the entire package is converted to a form that can be transmitted to the receiving device At the receiving machine. The message is unwrapped layer by layer with each process receiving and removing the dada meant for it.

