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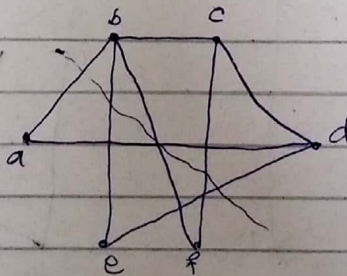
Question No 1

Ans:

Part A:

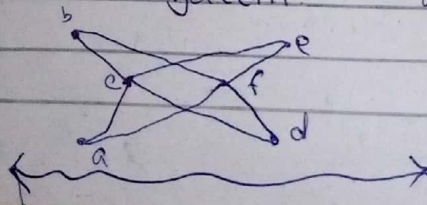
For the graphs given above, they are bipartite showing that they are 2-colorable.

The easy ones are the ones that are 2-colorable: you just find a 2-coloring i.e. a coloring of the vertices of the same color are never adjacent along an edge.

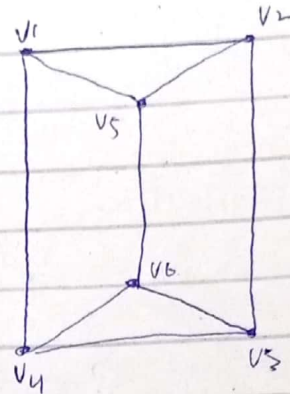
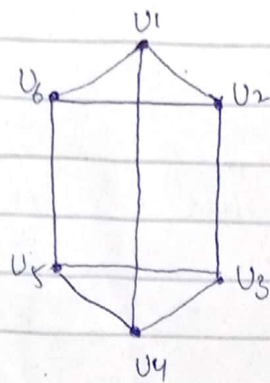


Part b: For the given graph above, they are bipartite showing they are 2-colorable

The easy ones are the ones that they are 2-colorable you just find a 2-coloring i.e. a coloring of the vertices of the same color are never adjacent along an edge.

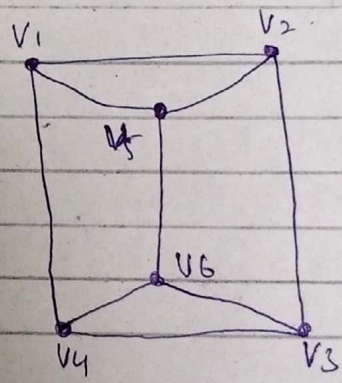
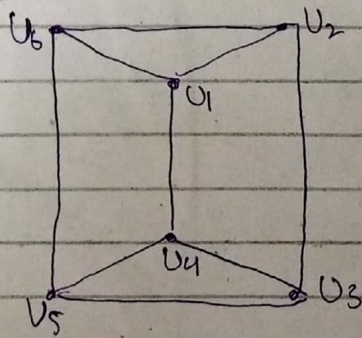


Question No: 2

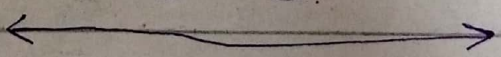


Answer: first we see the given graphs, they are isomorphic because when we keep vertex  $U_1$  and  $U_4$  under the square,  $U_1$  and  $V_5$  are same,  $U_4$  and  $V_6$  are same, therefore they are isomorphic

Example:



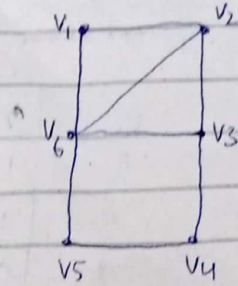
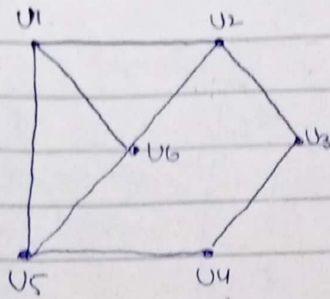
{ same }





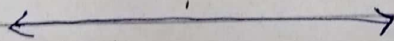
Question No. 2

Part 2:



Answer:

First we see the give graphs; they are non-isomorphic because when we compare vertices  $u$  and  $v$  they are not same, therefore they are not isomorphic.



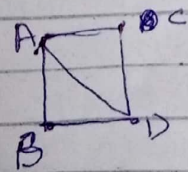
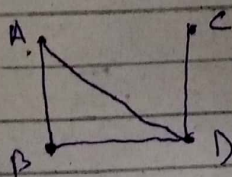
Question No. 3

Part 1:

	A	B	C	D
A	0	1	0	1
B	1	0	0	1
C	0	0	0	1
D	1	1	1	0

	A	B	C	D
A	0	1	1	1
B	1	0	0	1
C	1	0	0	1
D	1	1	1	0

Solution

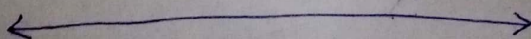


Adjacency lists:-

- $a = B, d$
- $b = A, d$
- $c = d$
- $d = A, B, C$

- $a = c, B, d$
- $b = A, d$
- $c = A, d$
- $d = A, B, C$

$\Rightarrow$  Not Isomorphic

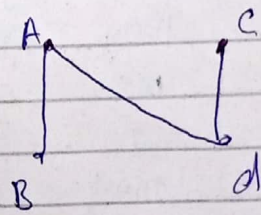
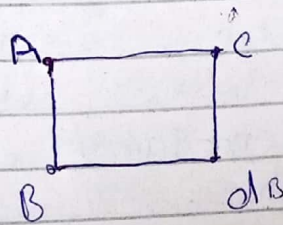


part:2

	A	B	c	d
A	0	1	1	0
B	1	0	0	1
c	1	0	0	1
d	0	1	1	0

	A	B	c	d
A	0	1	0	1
B	1	0	0	0
c	0	0	0	1
d	1	0	1	0

Solution:

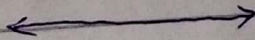


Adjacency lists:

A: c, B  
B: A, d  
c: A, d  
d: B, c

A: B, d  
B: A  
c: d  
d: A, c

Not Isomorphic :-



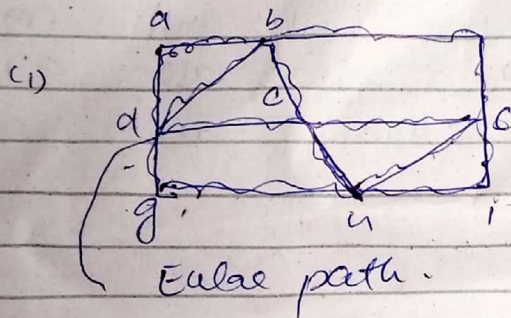


Q4

Answer

Euler circuit means a graph where we allow loops and parallel edges

and the path has the same initial and terminal vertices, we call it an Euler circuit.



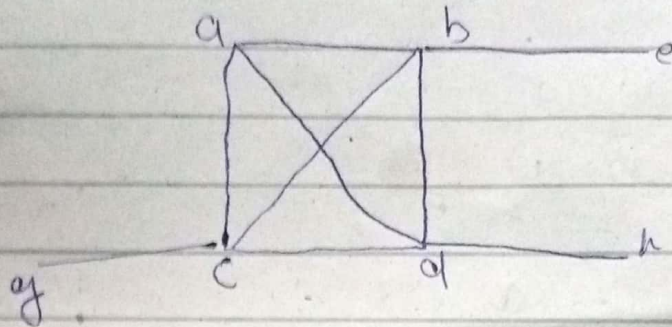
initial path terminal points is once

Start from a and end to a

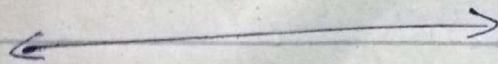
it is a Euler circuit.

Question No. 5. part: A

Solution.



The graph does not have a Hamiltonian path. Since there are two vertices of degree one, however if there are a Hamiltonian path then that are at 2 vertices of degree 1.





## Question no 15 Part B

### Solution

Let us determine the degree of every vertices in the given graph

$$\deg(a) = 2$$

$$\deg(b) = 4$$

$$\deg(c) = 2$$

$$\deg(d) = 3$$

$$\deg(e) = 3$$

We then note that Dirac's theorem is not satisfied (Since some degrees are less than  $n/2 = 5/2 = 2.5$ ), but this does not necessarily mean that no Hamilton circuit exists

However, we do note that given graph contain the cycle  $C_4$  and the cycle  $C_3$  with in the given graph forms a Hamilton circuit (as the circuit) will pass through all vertices exactly ones).

A possible Hamilton circuit is thus the path of  $C_4 = a, b, c, d, e, a$

