

1

NAME: YASEEN-KHAN B.S (SE)A 14461

Q1:) In a block of address, we know the IP address of one host is $101.10.11.X/ID_{4+5}$. What are first address (network address) and the last address (limited broadcast address) in this block.

Given:- IP of one host = $101.10.11.X/ID_{4+5}$

Required:- $X = ?$, $ID_{4+5} = ?$

First address (network address) = ?

Last address (limited broadcast address) = ?

Solution:-

$$101.10.11.X/ID_{4+5} \quad \text{--- (1)}$$

$$X = \text{Sum of ID}$$

$$ID = 14461 \Rightarrow \text{Sum} = 1+4+4+6+1 = 16$$

$$X = 16 \quad \text{--- (2)}$$

$$ID_{4+5} = \text{Sum of 4th and 5th digit of ID}$$

$$= 4^{\text{th}} = 6, \quad 5^{\text{th}} = 1$$

$$= 6+1$$

$$ID_{4+5} = 7 \quad \text{--- (3)}$$

NAME: YASEEN-KHAN

BS (SE) A

14461

Putting values of (2) and (3) in (1)

$$IP = 101.10.11.16/7$$

(a) First address (Network address)

$$IP = 101.10.11.16/7$$

Converting to Binary Notation.

$$101 = 01100101$$

$$10 = 00001010$$

$$11 = 00001011$$

$$16 = 00010000$$

In Binary Notation

$$01100101.00001010.00001011.00010000$$

⇒ First address can be found by:

$$32 - n \Rightarrow n = \text{net mask} = 7$$

$$32 - 7 = 25$$

We will convert 25 right most bit's to 0's for first address

$$01100101.00001010.00001011.00010000$$

Converting right most 25 bits to 0's

NAME: YASEEN-KHAN

BS (SE) A

14461

01100100 · 00000000 · 00000000 · 00000000

Converting above to Decimal notation

$$\begin{aligned}
 01100100 &= 0 \times 2^7 + 1 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 0 \times 2^0 \\
 &= 64 + 32 + 4 \\
 &= 100
 \end{aligned}$$

$$\begin{aligned}
 00000000 &= 0 \times 2^7 + 0 \times 2^6 + 0 \times 2^5 + 0 \times 2^4 + 0 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 \\
 &\quad + 0 \times 2^0 \\
 &= 0
 \end{aligned}$$

we get

100 · 0 · 0 · 0 / 7 as first address

b) Last address.

$$IP = 101 \cdot 10 \cdot 11 \cdot 16 / 7$$

Converting to Binary Notation

$$101 = 01100101$$

$$10 = 00001010$$

$$11 = 00001011$$

$$16 = 00010000$$

In - Binary Notation

$$01100101 \cdot 00001010 \cdot 00001011 \cdot 00010000$$

YASEEN-KHAN

BS(SE)A

14461

⇒ Last address can be found by changing $32-n$ bits of most right to '1'.

$$32-n \Rightarrow 32-7 \Rightarrow 25$$

∴ netmask n

Converting 25 right most bits to '1'

01100101.00001010.00001011.00010000
 converting these bits to 1

01100101.11111111.11111111.11111111

Converting above to Decimal Notation

$$\begin{aligned} 01100101 &= 0 \times 2^7 + 1 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= 64 + 32 + 4 + 1 \\ &= 101 \end{aligned}$$

$$\begin{aligned} 11111111 &= 1 \times 2^7 + 1 \times 2^6 + 1 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 \\ &= 128 + 64 + 32 + 16 + 8 + 4 + 2 + 1 \\ &= 255 \end{aligned}$$

$$101.255.255.255 / 7$$

(Last address)

NAME: YASEEN-KHAN BS (SE) A 14461

Q2) Take your Roll No as decimal Notation now convert it into Binary Notation. Draw the graph of the NRZ-L scheme using the Binary Notation of your Roll no as data stream, assuming that last signal level has been positive.

ID = 14461

14461
Converting to ~~Hex~~ Binary notation

$$14461 = 11100001111101$$

Now creating NRZ-L scheme.

As we have to keep our last value positive, our last value is already positive, No changes will be made.

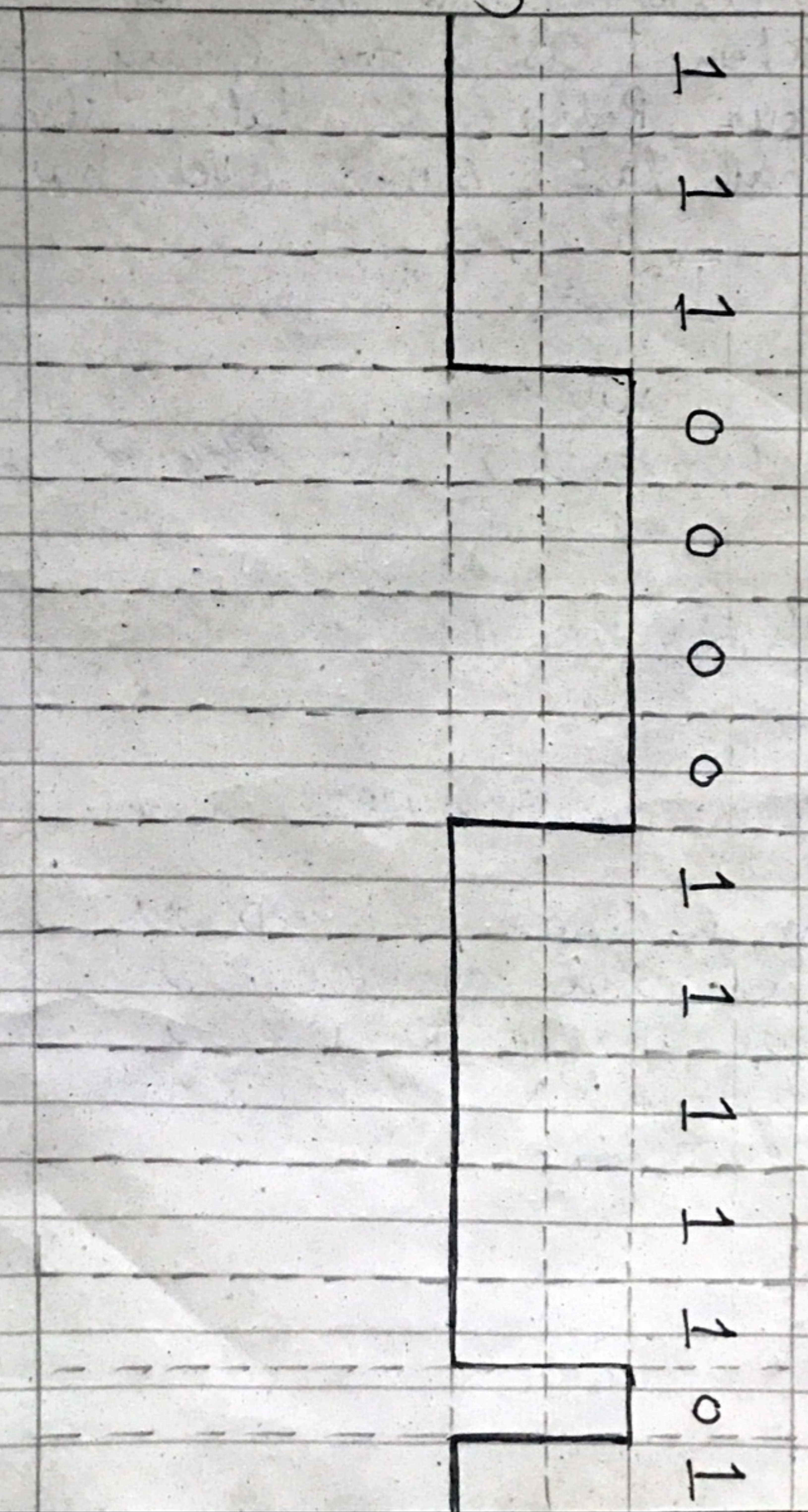
2	14461
2	7230-1
2	3615-0
2	1807-1
2	903-1
2	451-1
2	225-1
2	112-1
2	56-0
2	28-0
2	14-0
2	7-0
2	3-1
	1-1

YASEEN - KHAN

BS (SE) A

14421

NRZ-L
(Scheme)



YASEEN-KHAN

7

BS (SE) A

14461

Q3)

Window Size = ID_{last}

$ID = 14461$

$ID_{last} = 1$

Window size = 1

(a) Before A sends any frame.

Sender:-

0 1 2 3 4 5 6

windows of PDU transmitted = 1 bit

Receiver:-

0 1 2 3 4 5 6

window

(b) After A sends frame 0, 1, 2 ---

Sender:-

0 1 2 3 4 5 6

A shrunk it's window, keeps copy of PDU

Receiver:-

0 1 2 3 4 5 6

Receives all data.

14461

Q4)

$$\frac{160 \cdot (x) \cdot (10_{3+4}) \cdot 0}{16}$$

$$\frac{160 \cdot 16 \cdot 10 \cdot 0}{16}$$

(a)

16 customers needs 64 addresses

$$6 (\log_2 64)$$

$$36 \text{ bits for each host.}$$

$$\text{length of prefix}$$

$$32 - 6 = 26$$

1st customer:-

$$160 \cdot 16 \cdot 10 \cdot 0 / 26$$

$$160 \cdot 16 \cdot 10 \cdot 63 / 26$$

2nd customer:-

$$160 \cdot 16 \cdot 11 \cdot 0 / 26$$

$$160 \cdot 16 \cdot 11 \cdot 127 / 26$$

3rd customer:-

$$160 \cdot 16 \cdot 12 \cdot 0 / 26$$

$$160 \cdot 16 \cdot 12 \cdot 191 / 26$$

}

16th Customer:-

$$160 \cdot 16 \cdot 25 \cdot 0 / 26$$

$$160 \cdot 16 \cdot 25 \cdot 255 / 26$$

Total = $16 \times 64 = 1024$ addresses.

Yascan - Khan

BS(SEE)A

14461

(b) Second group has 64 customers
each needs 32 addresses
 $5 (\log_2 32)$
25 bits for each host.

length of prefix:

$$32 - 5 = 27$$

1st customer:-

$$160.16.26.0/27 \quad 160.16.26.31/27$$

2nd customer:-

$$160.16.27.32/27 \quad 160.16.27.63/27$$

3rd customer:-

$$160.16.28.64/27 \quad 160.16.27.95/27$$

}

64th customer:-

$$160.16.89.255/27 \quad 160.16.89.255/27$$

$$\text{Total addresses} = 64 \times 32 = 2048$$

Yaseen-Khan

BSC(SE)A

14461

(c) Third group has 64 customers each needs 16 addresses.

$i = 6$ ($\log_2 16$)
24 bits for each host.

$$32 - 4 = 28$$

1st customer:-

$$160.16.90.0/28$$

$$160.16.90.15/28$$

2nd customer:-

$$160.16.91.16/28$$

$$160.16.91.31/28$$

3rd Customer:-

$$160.16.92.32/28$$

$$160.16.92.47/28$$

}

164th customer:-

$$160.16.153.255/28$$

$$160.16.153.255/28$$

$$\text{Total} = 64 \times 16 = 1024 \text{ address}$$

Yaseen-Khan

BS(SE) A

(4461)

Numbers of granted address = 65,536
addresses used = 40,96
Available addresses = 61,440