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Subject: Compiler Construction

EVEN ROLL NUMBERS

Q1. Construct regular expression defining each of the following language over the alphabet $\Sigma=\{a, b\}$.

i. All words having even length

Ans: EVEN-EVEN = $\{\Lambda, aa, bb, aaaa, aabb, abab, abba, baab, baba, bbaa, bbbb, \dots\}$,

its regular expression can be written as

$(aa+bb+(ab+ba)(aa+bb)^*(ab+ba))^*$ ii. **All words**

having at least three a and three b

Ans: $r1=(aaaa, baaa, abbb, bbbb)$ Its regular

expression can be written as $r1=(a+b)^*(aaa+bbb)$

$r2=(a+b)^*aaa+(a+b)^*bbb$ iii. **All words having at**

least double a or triple b

Ans: $r1(aaa, baa, abaa, baaa)$

Its regular expression can be written as

$r1=(aa+ba)^*a+(aba+baa)^*a$

iv. All words starts with four a or triple b.

Ans: $r1(bbba, bbbab, bbbba, bbbb)$ Its

regular expression can be written as

$r1=b^*(bba+bbab)+b^*(bbba+bbb)$

Q2. For figure 3 if q_0 is the initial state, the draw a transition table for it.

Ans: The states are named by capital letters this time for a bit of variation: $Q = \{q_0, q_1, q_2, q_3, q_4\}$. While it is common to use symbols $q_i, i \in \mathbb{N}$ to name states, we can pick any names we like. Another common choice is to use natural numbers; i.e., $Q \subset \mathbb{N} \wedge Q$ is finite.

The representation of the above DFA as a transition table is:

Q	0	1
→ q ₀	q ₁	q ₂
q ₁	q ₃	q ₂
q ₂	q ₁	q ₄
* q ₃	q ₃	q ₂
* q ₄	q ₁	q ₄

Q3. Define what is Finite Automaton. What can be the regular expression of the diagram given in figure 1.

Ans: A finite automaton (FA) is a simple idealized machine used to recognize patterns within input taken from some character set (or alphabet) C . The job of an FA is to accept or reject an input depending on whether the pattern defined by the FA occurs in the input

A Finite automaton (FA), is a collection of the followings

- 1) Finite number of states.
- 2) Finite set of input letters (Σ) from which input strings are formed.
- 3) Finite set of transitions.

Transition table

State	Reading a	Reading b
x-	z	y
y	y	y
Z+	z	z

$\Sigma = \{a, b\}$, starting with a.

regular expression = $a(a + b)^*$

Q4. Draw a transition table for the diagram given in figure 2. (0) is the starting state Ans:

state	a	b	c	d	e	f	g	h
→ 0	1	0	0	0	2	0	0	0
1	1	3	0	0	2	0	0	0
2	1	0	0	0	2	4	0	0

3	5	0	6	0	7	0	0	0
4	8	0	0	0	9	0	10	0
5	5	0	6	0	7	0	0	0
6	5	0	6	11	7	0	0	0
7	5	0	6	0	7	12	0	0
8	8	12	0	0	9	0	10	0
9	8	0	0	0	9	0	10	0
10	8	0	0	0	9	0	10	13
11	5	0	6	0	7	0	0	0
12	14	0	15	0	16	0	17	0
*13	8	0	0	0	9	0	10	0
14	14	0	15	0	16	0	17	0
15	14	0	15	18	16	0	17	0
16	14	0	15	0	16	0	17	0
17	14	0	15	0	16	0	17	19
*18	14	0	15	0	16	0	17	0
*19	14	0	15	0	16	0	17	0

The initial state is identified by putting an arrow \rightarrow to the left of it, and all final states are similarly identified by a star $*$.